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(54) **SYSTEM AND METHOD FOR  
IMPLEMENTING PRIVACY CLASSES  
WITHIN AN AUTHENTICATION  
FRAMEWORK**

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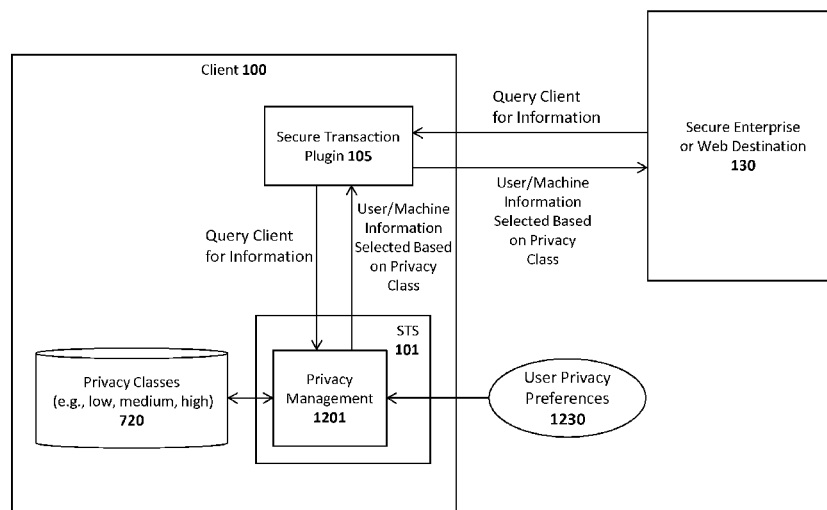
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(57) **ABSTRACT**

A system, apparatus, method, and machine readable medium are described for implementing privacy classes within an authentication framework. For example, one embodiment of a method comprises: transmitting a query for client information from a server to a client, the client information including information related to authentication devices coupled to the client; analyzing the query to determine an appropriate privacy class to be used for providing client information to the server; providing a subset of client information selected based on the determined privacy class, the subset of client information including the information related to the authentication devices coupled to the client; and using the subset of client information within an authentication framework to provide user authentication services over a network.

**33 Claims, 20 Drawing Sheets**



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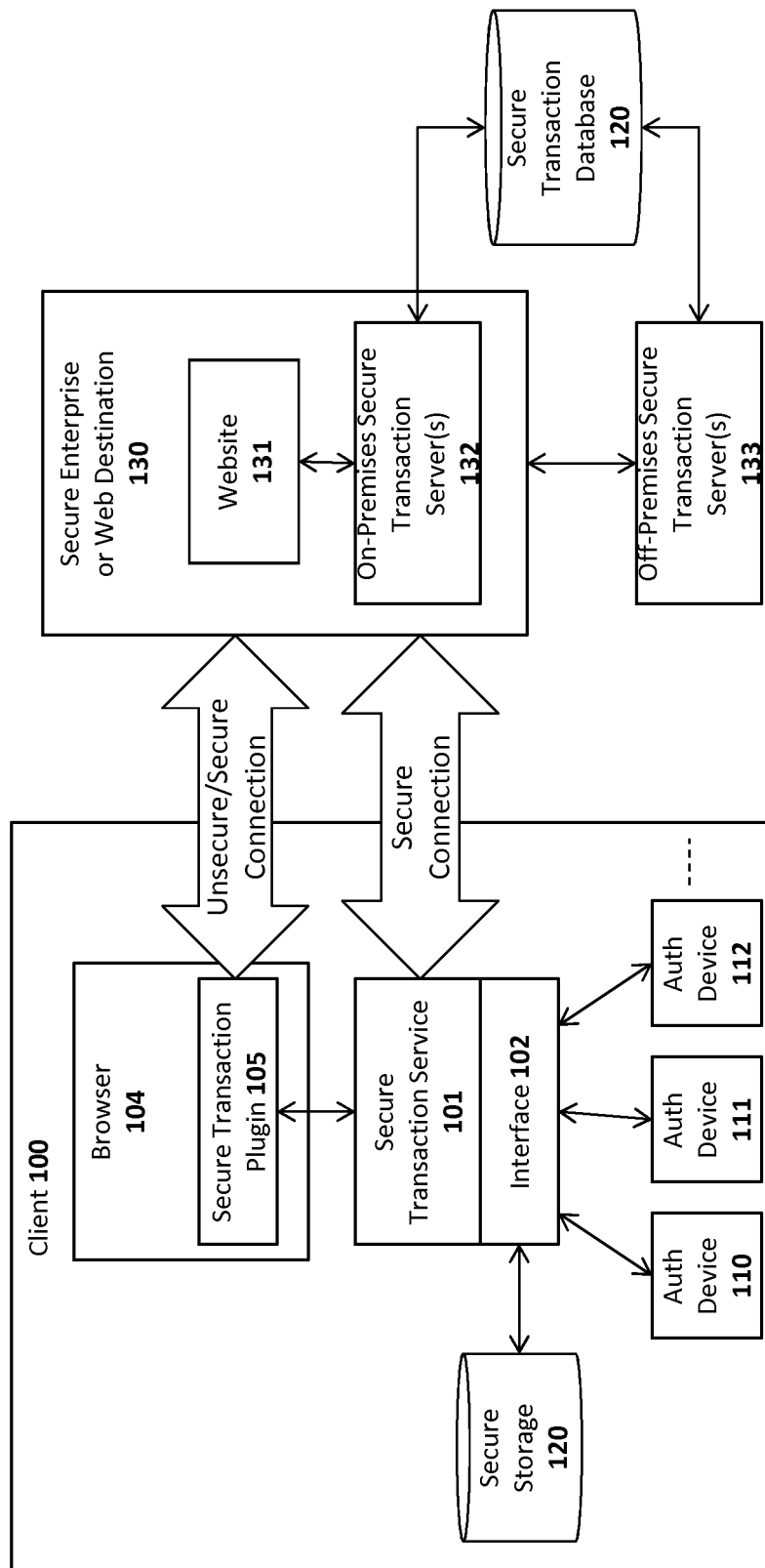


FIG. 1A

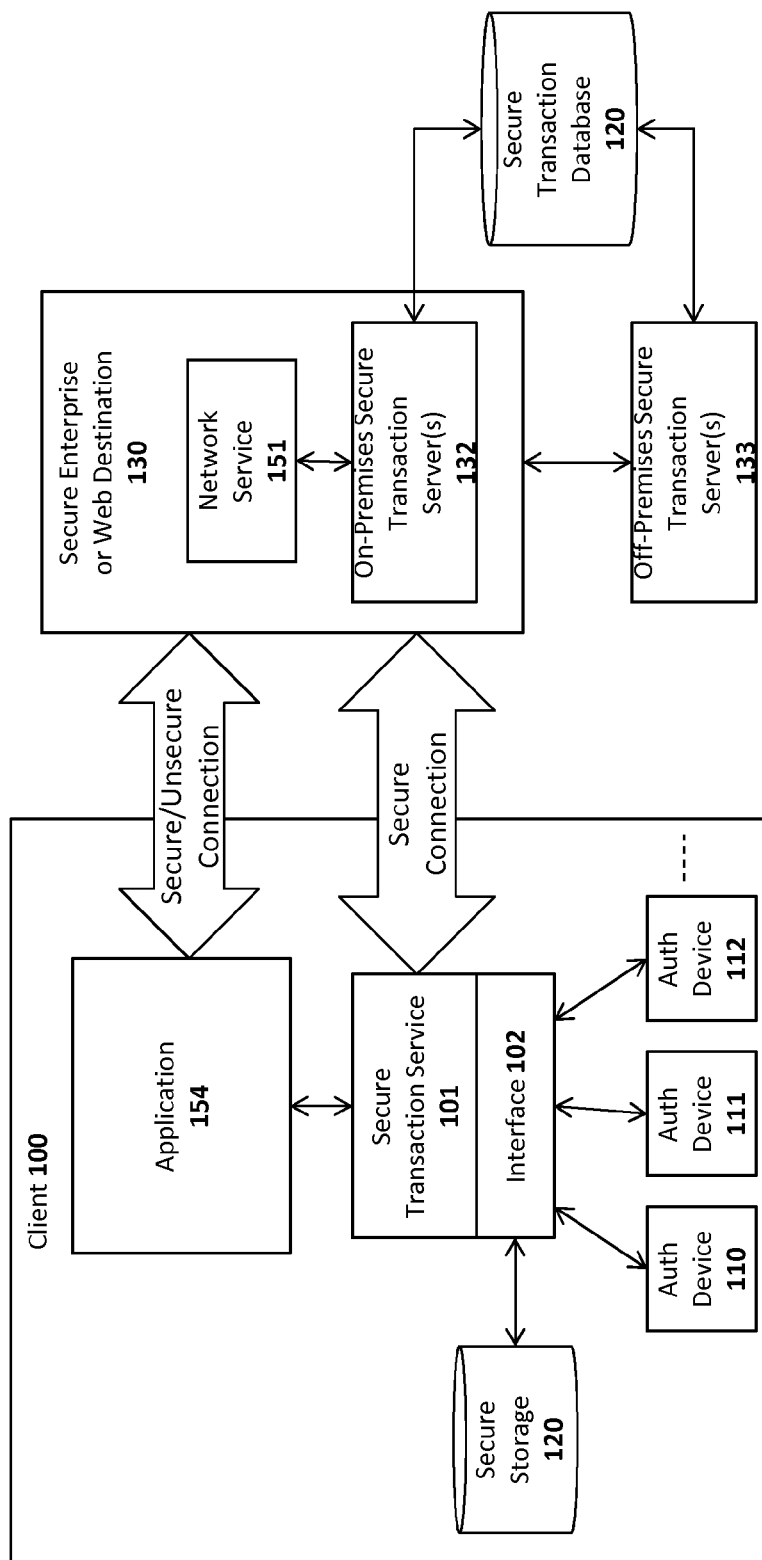
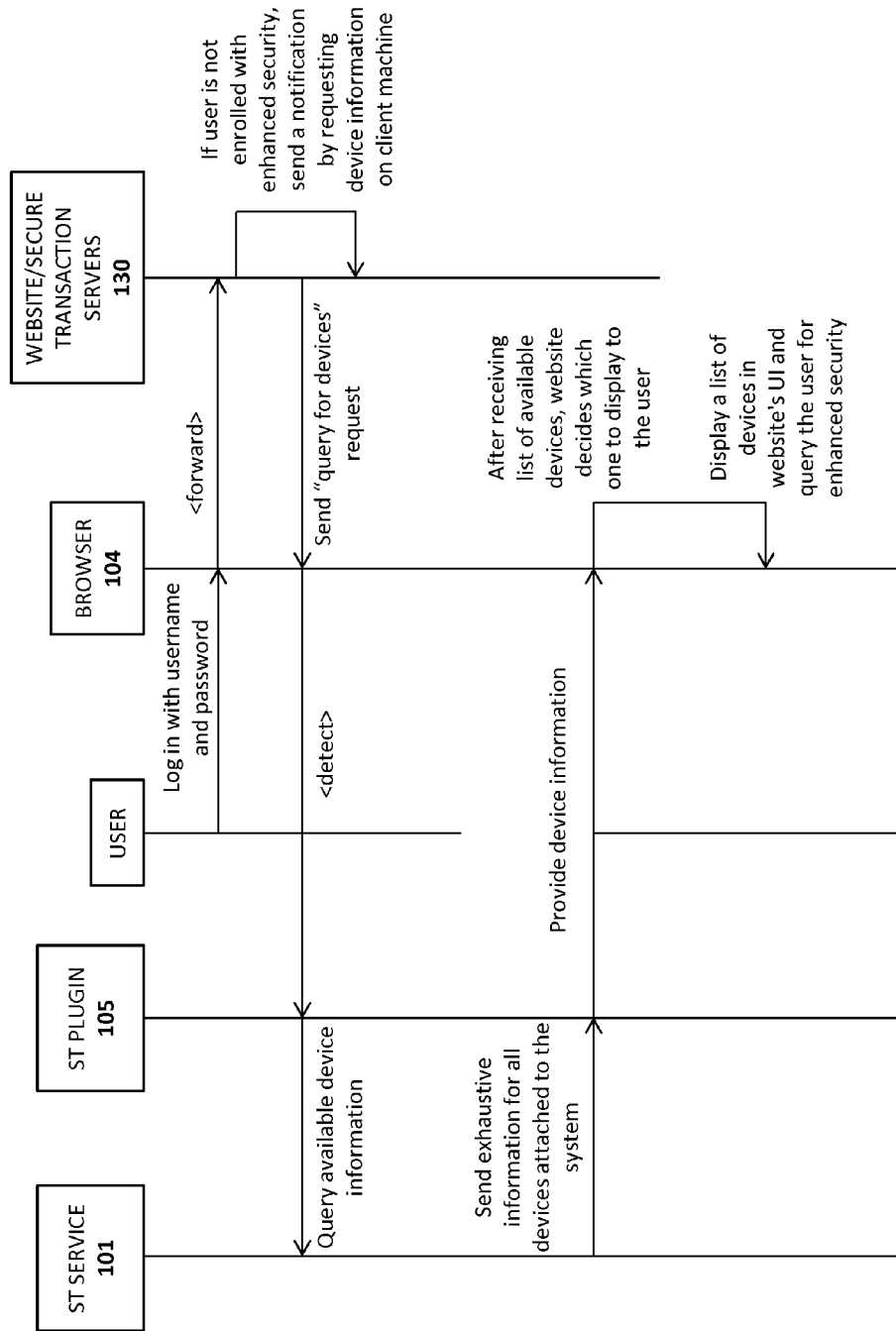


FIG. 1B



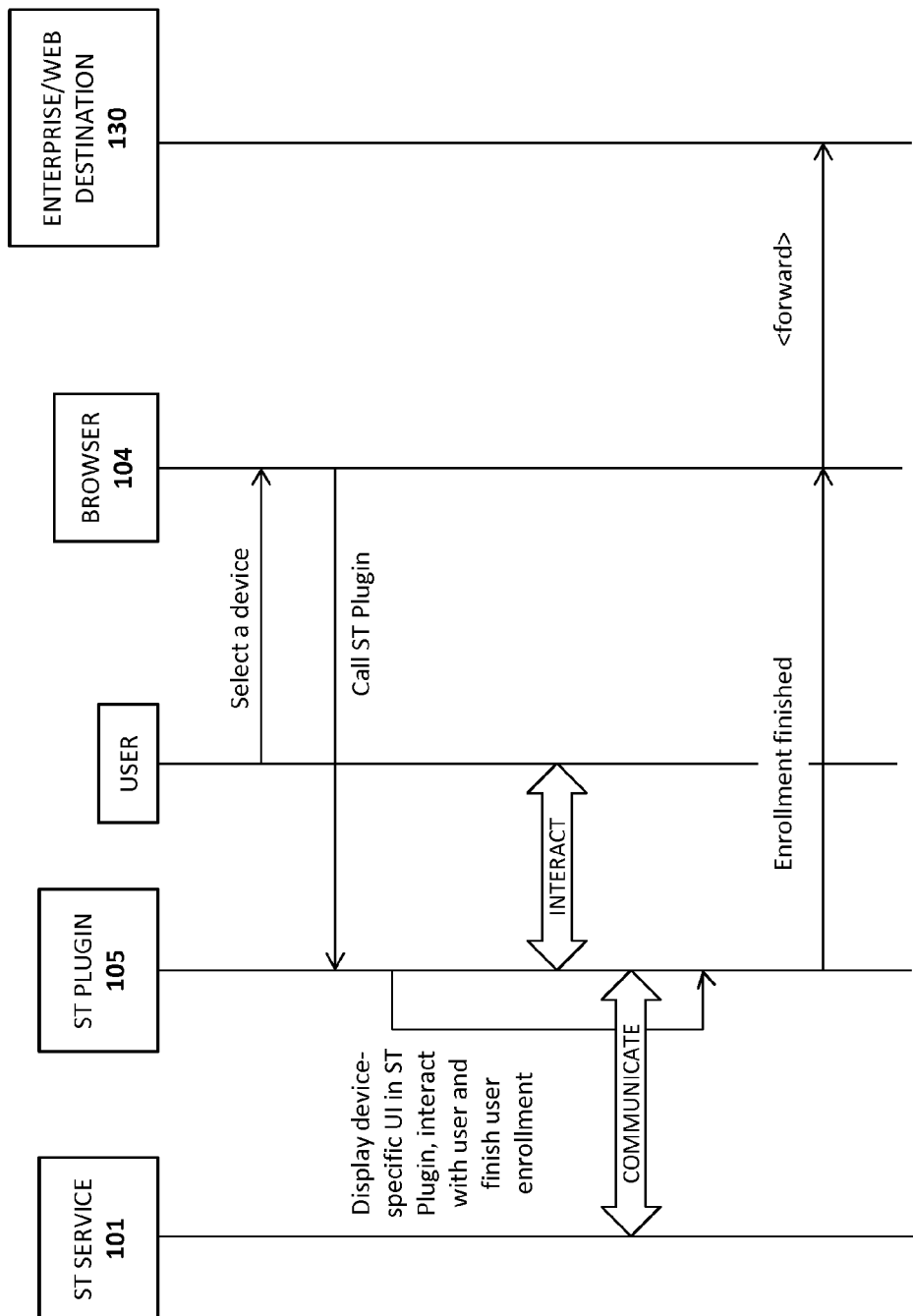


FIG. 3

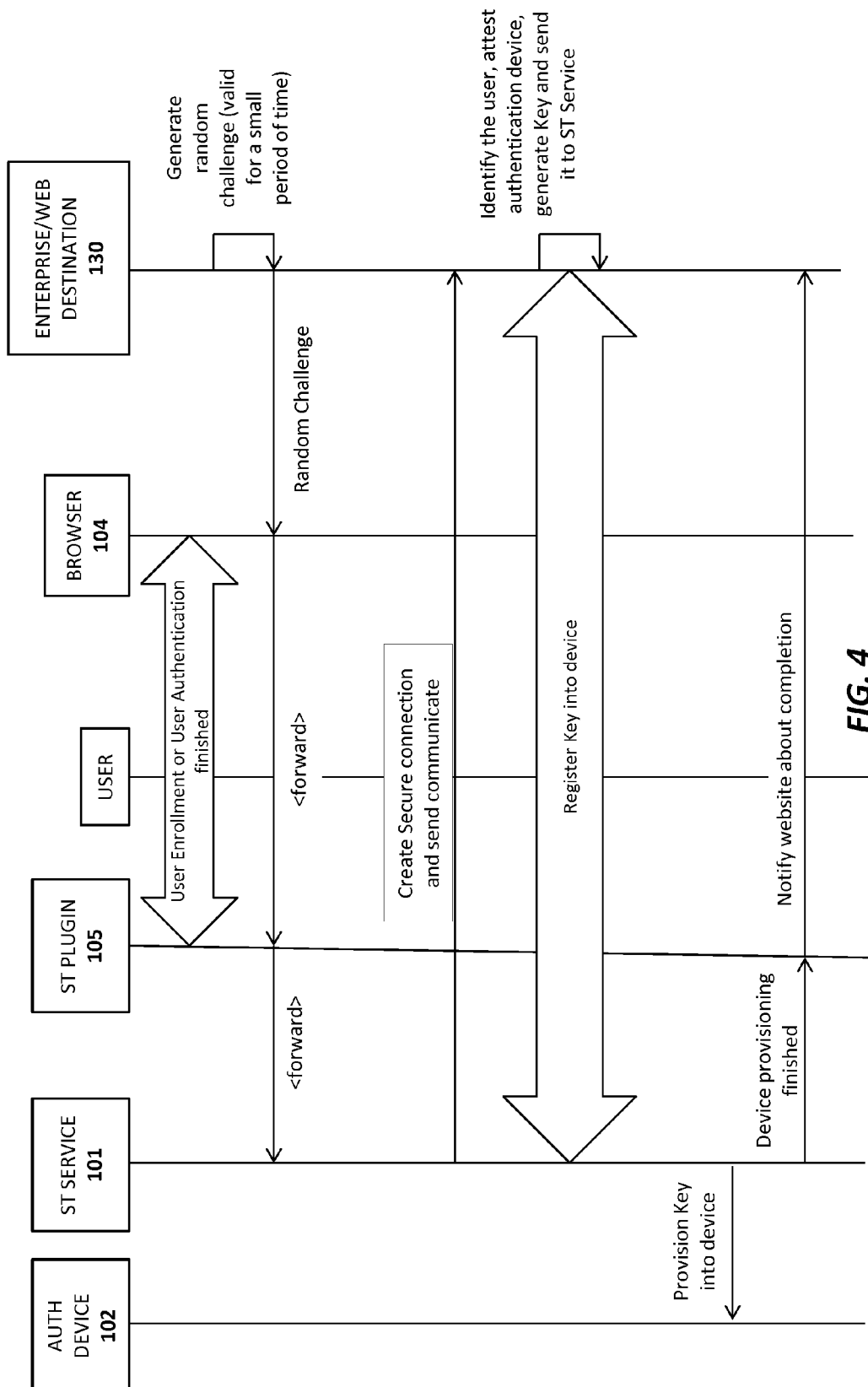
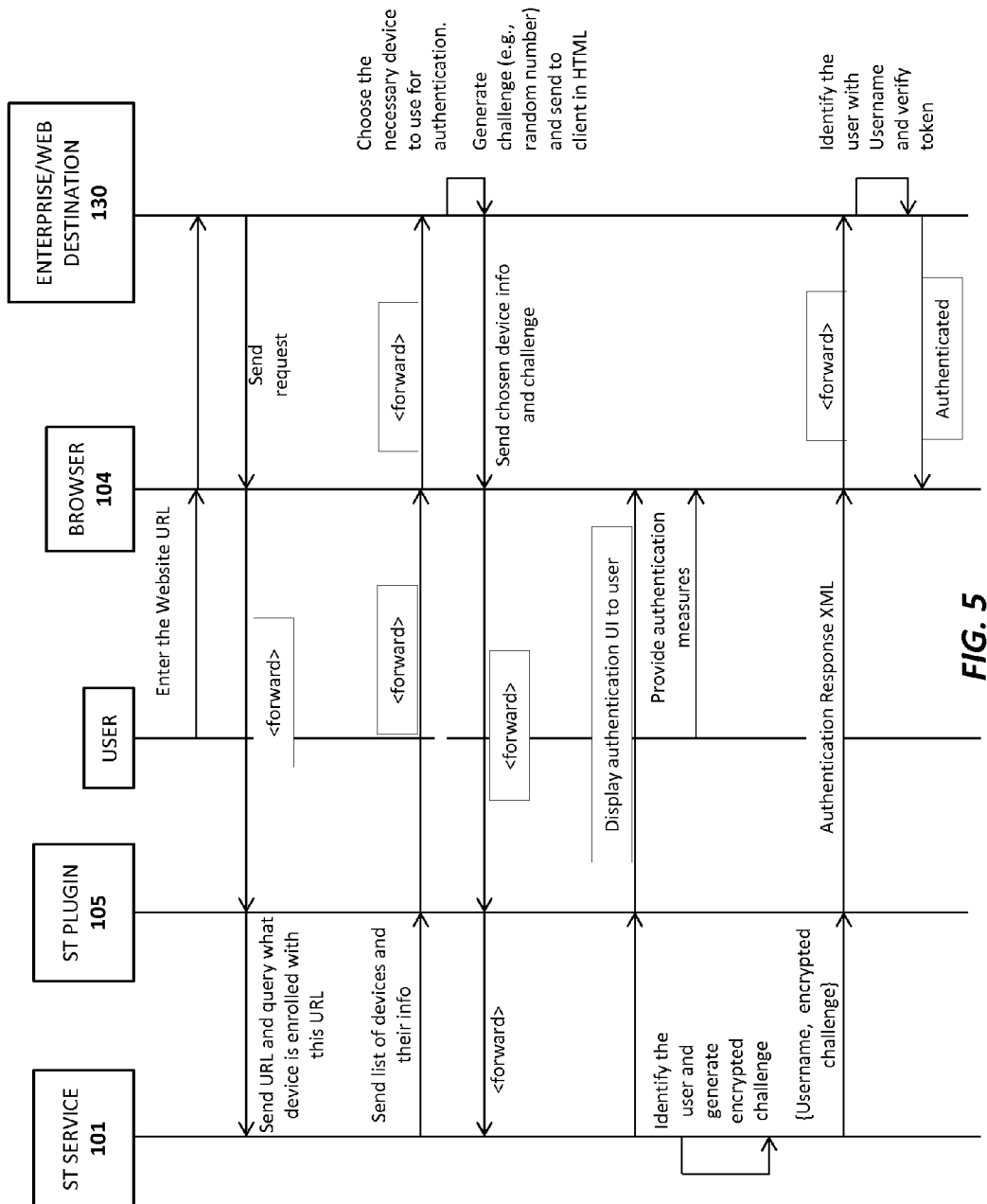


FIG. 4





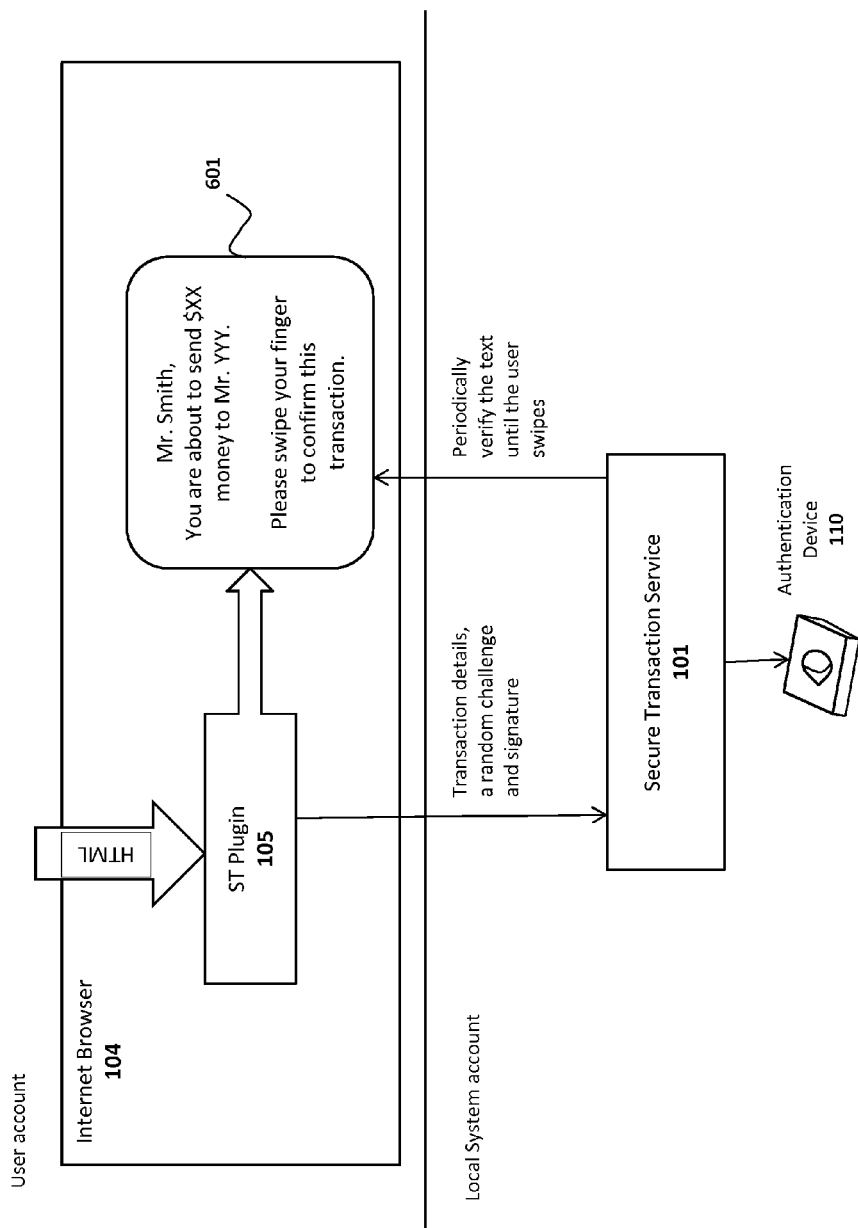


FIG. 6

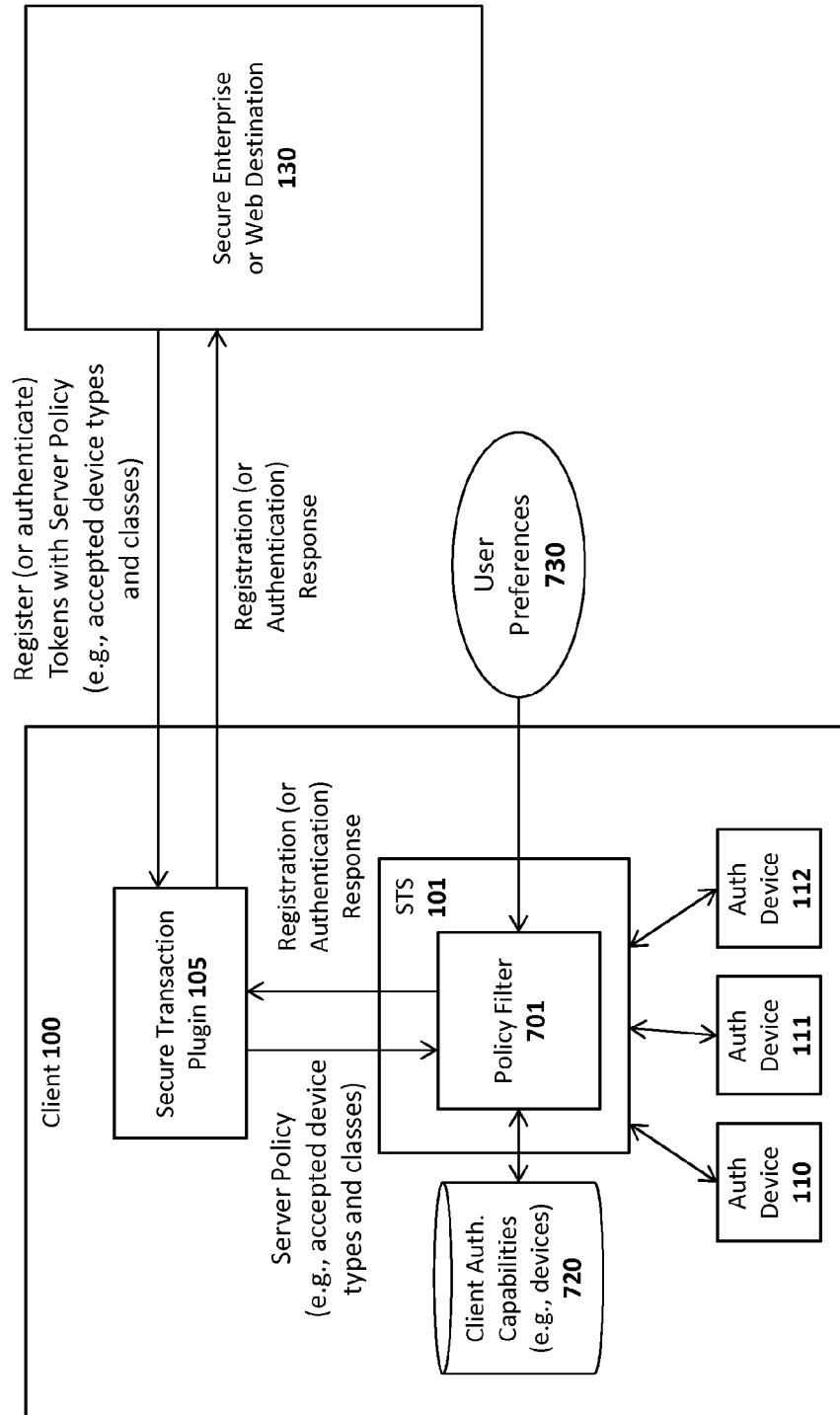
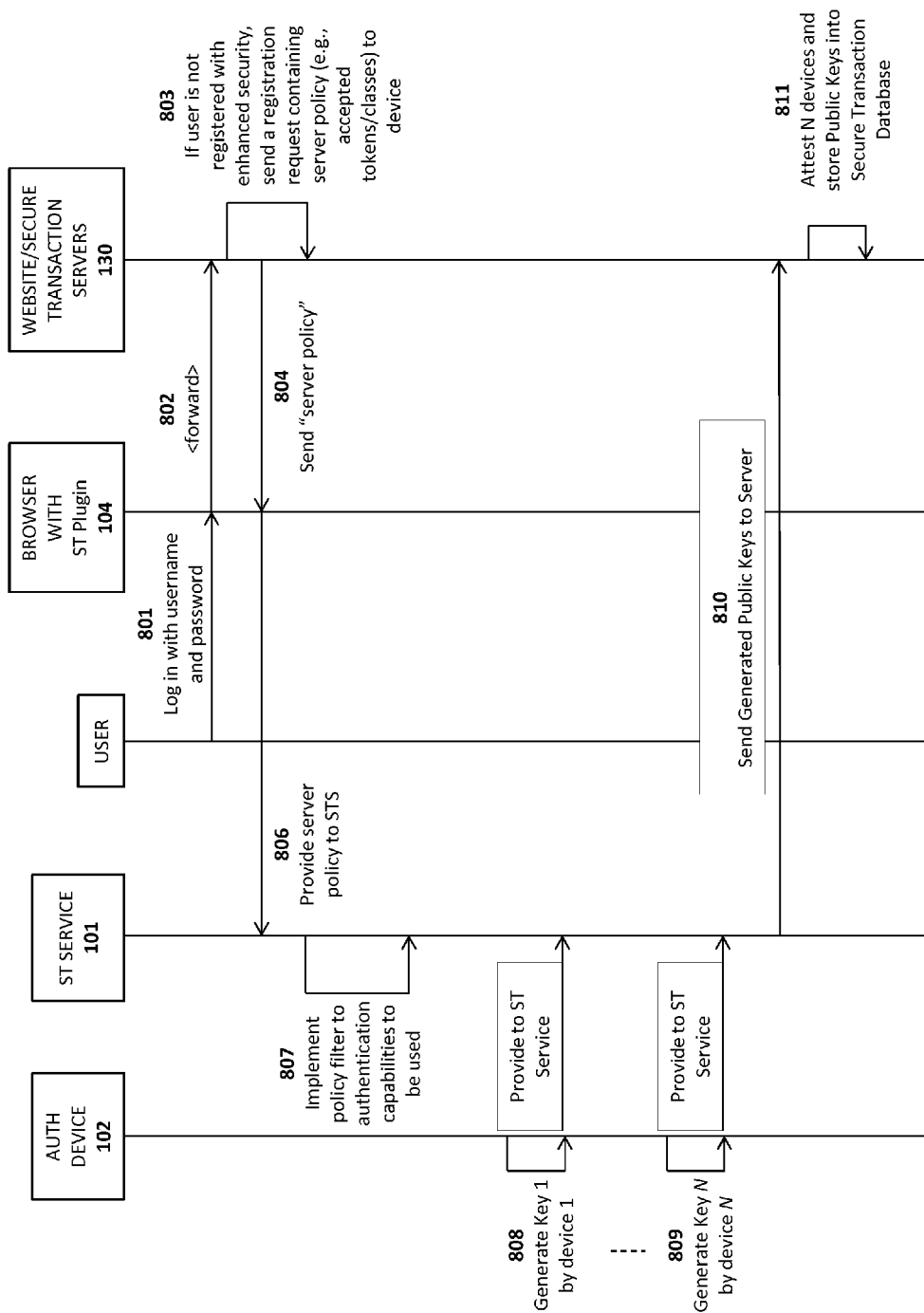


FIG. 7



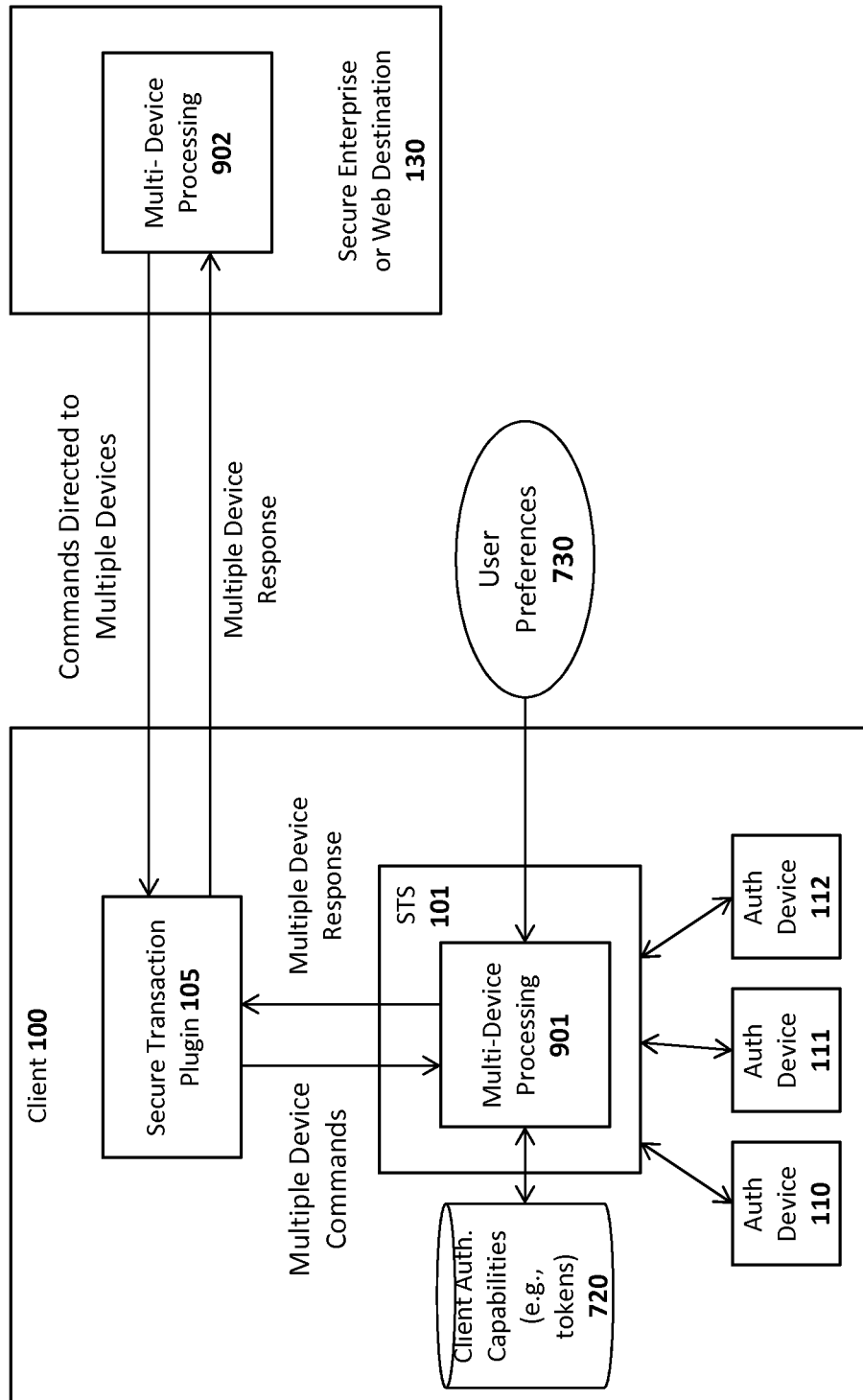


FIG. 9

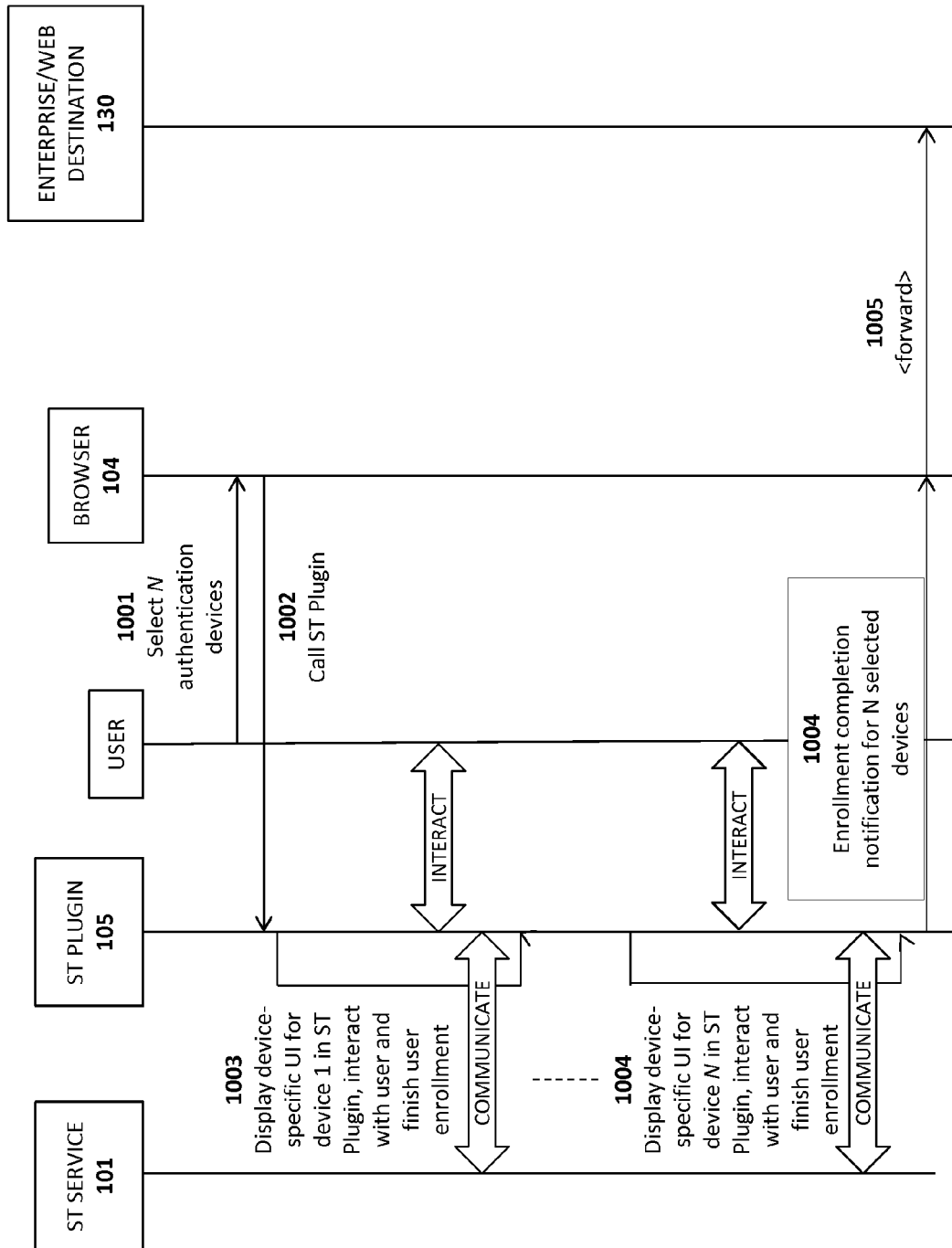


FIG. 10A

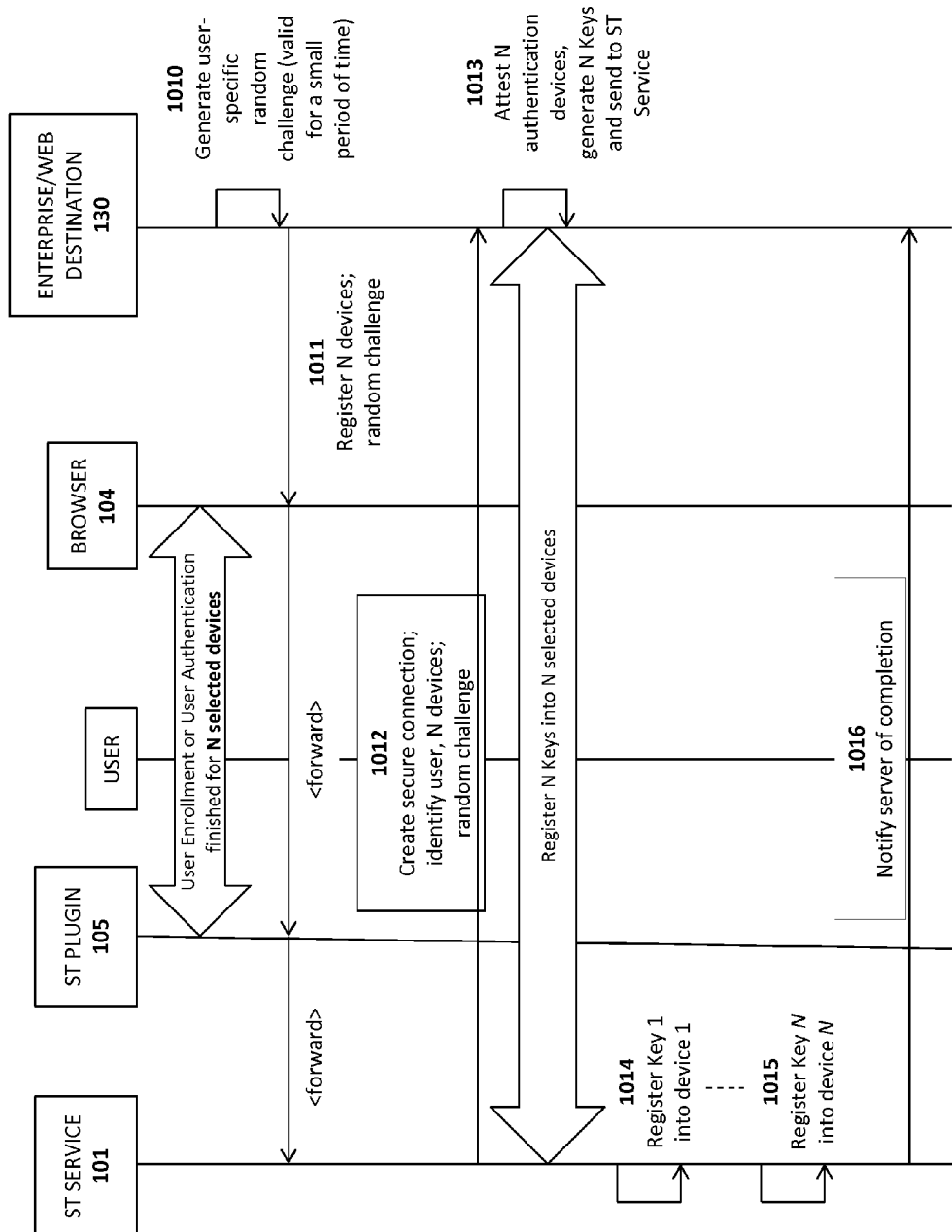


FIG. 10B

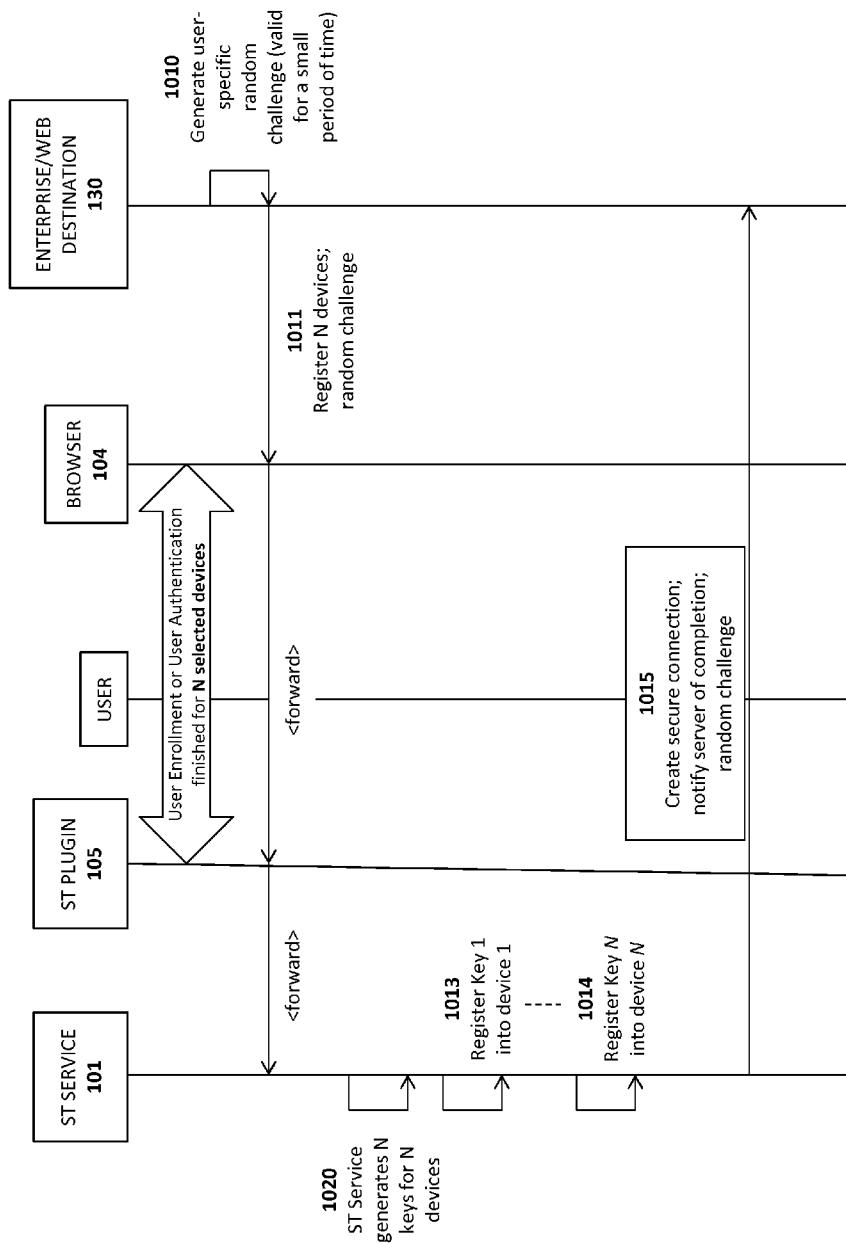


FIG. 10C

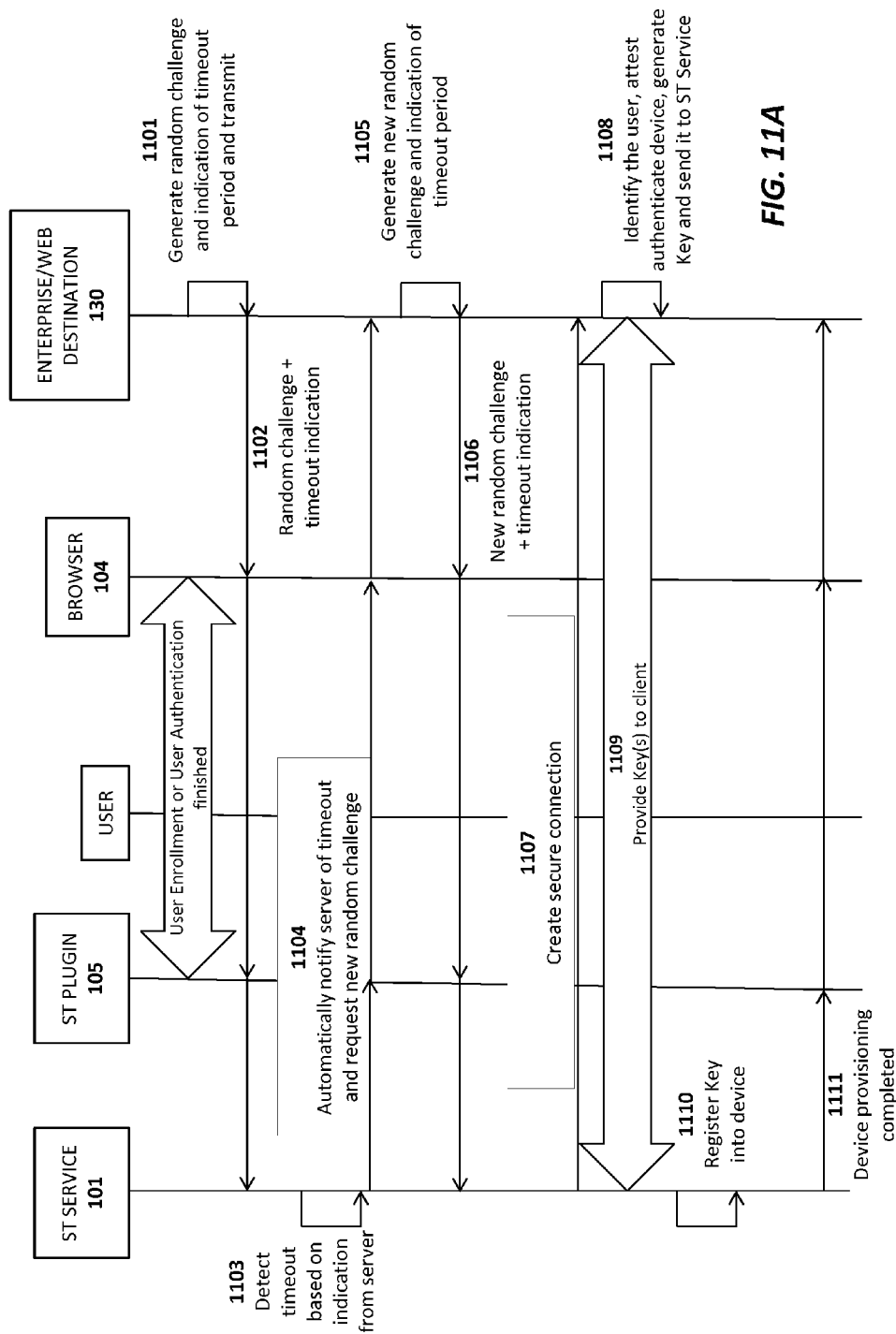
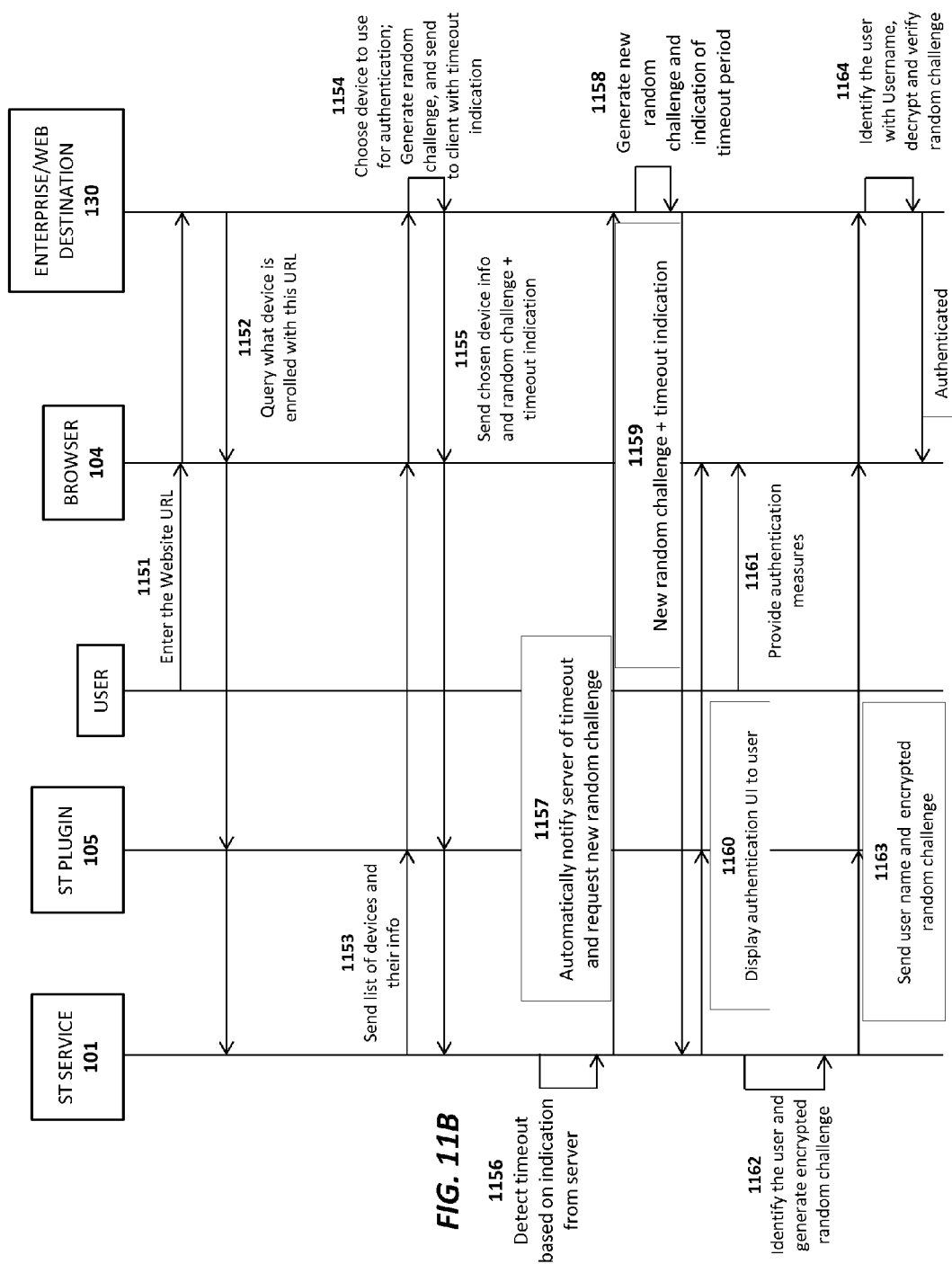


FIG. 11A





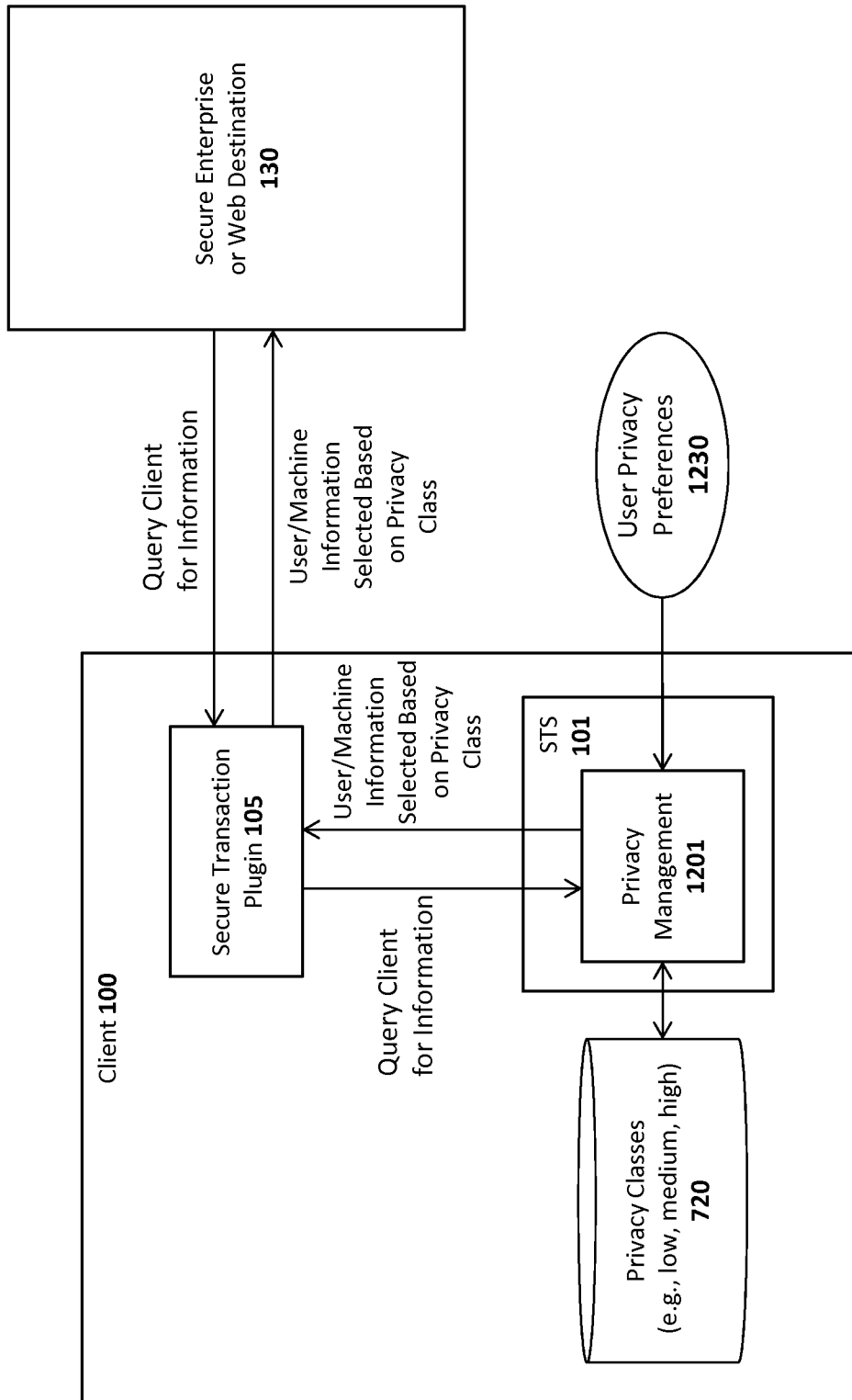


FIG. 12

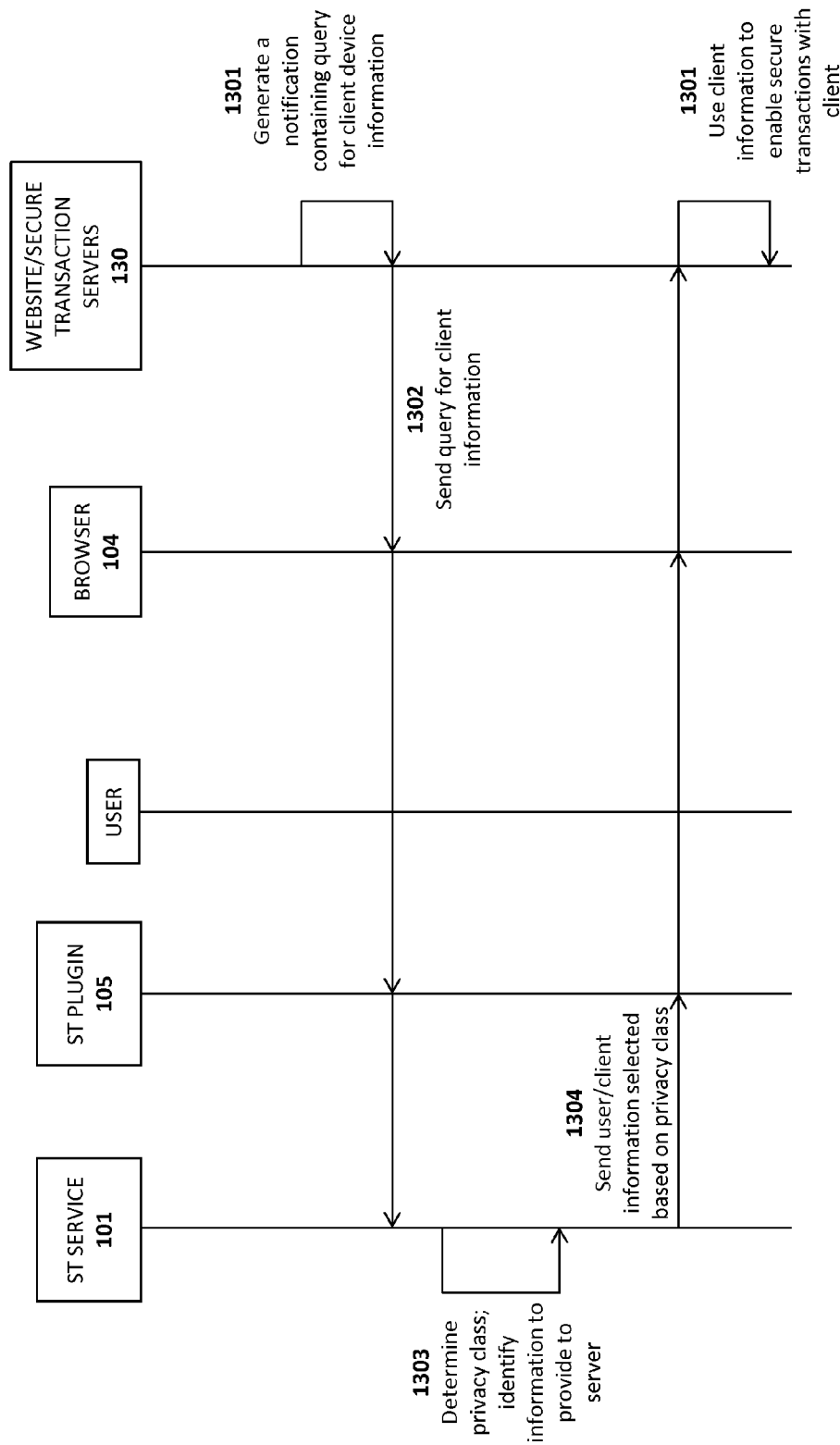


FIG. 13

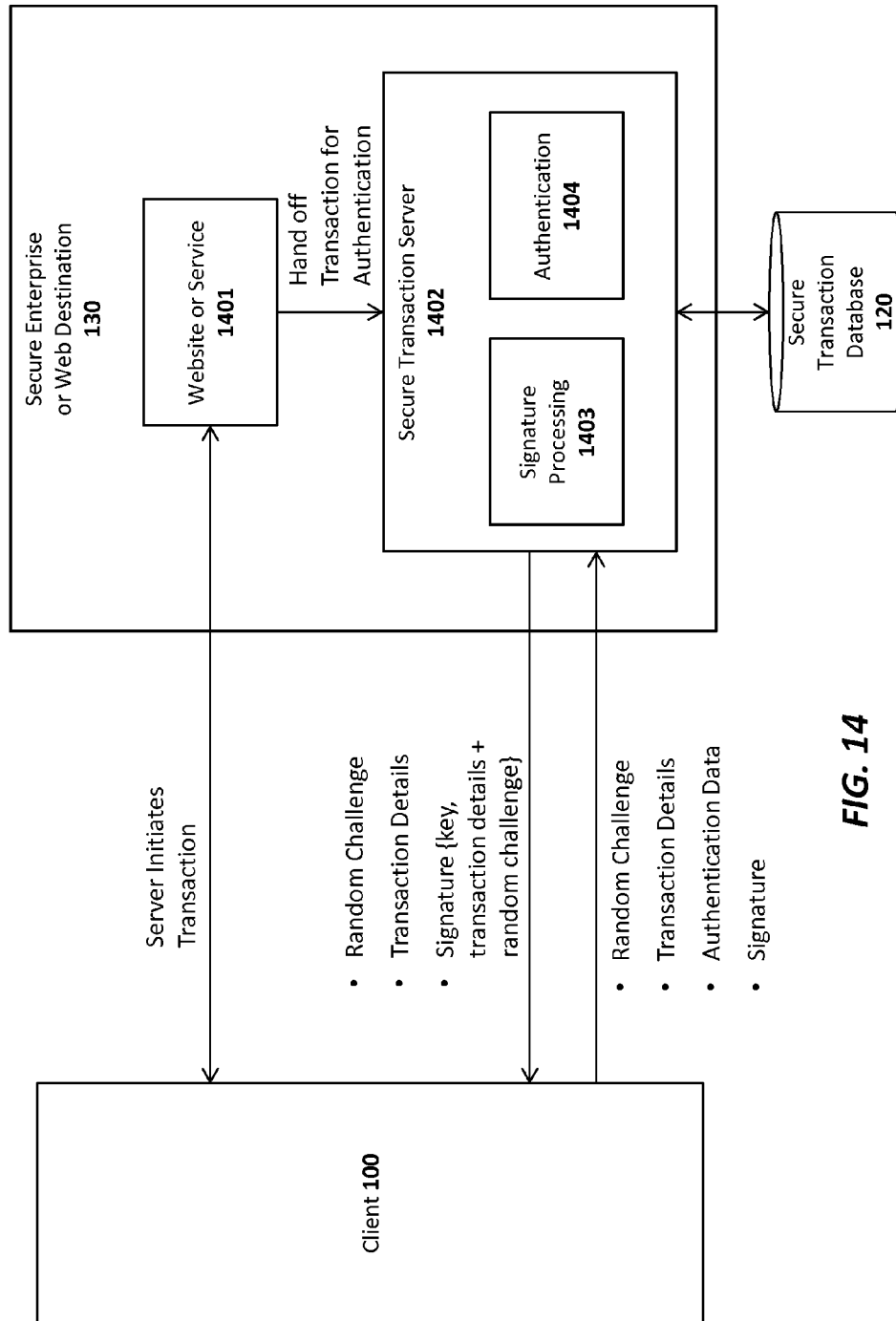
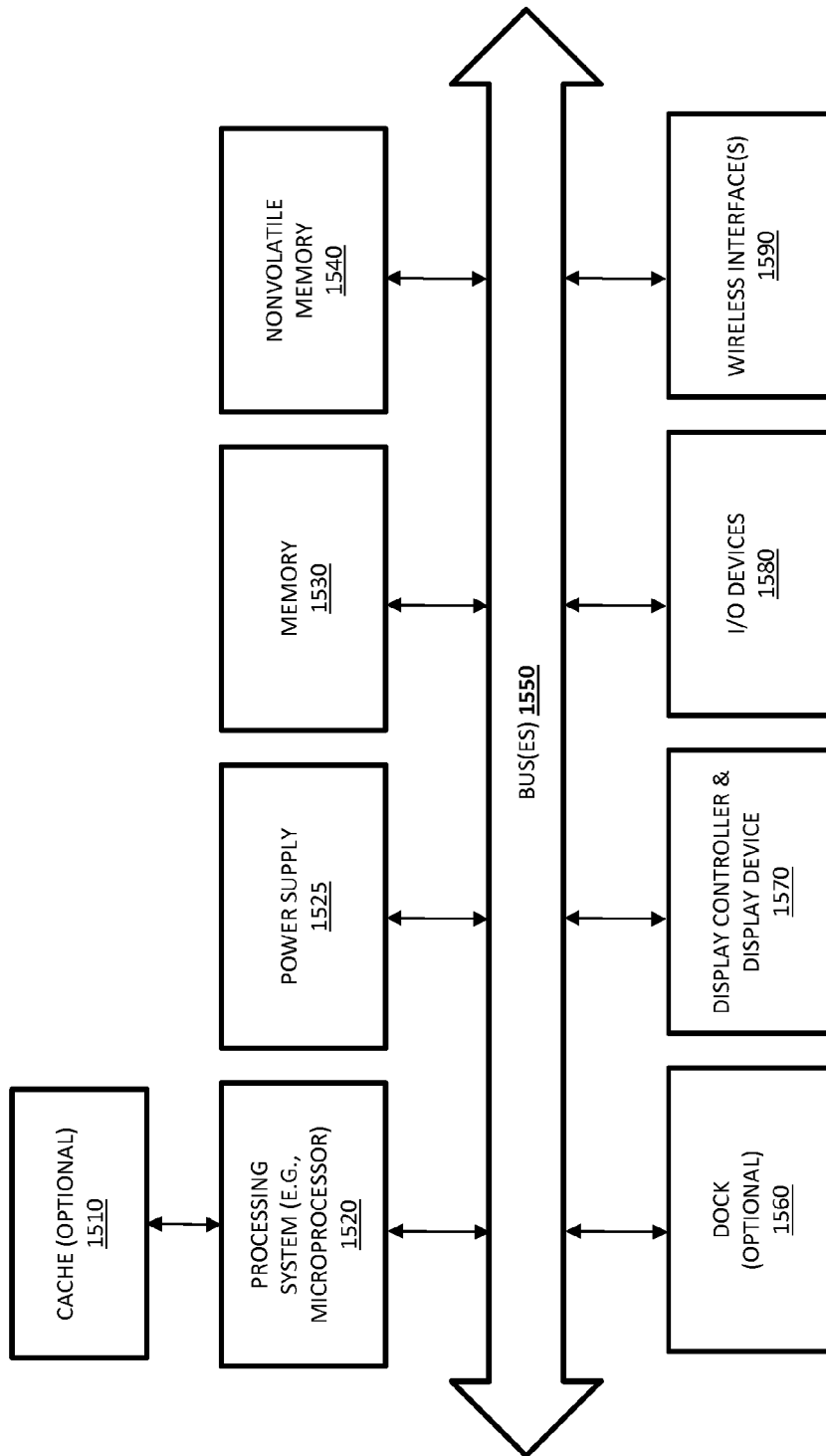
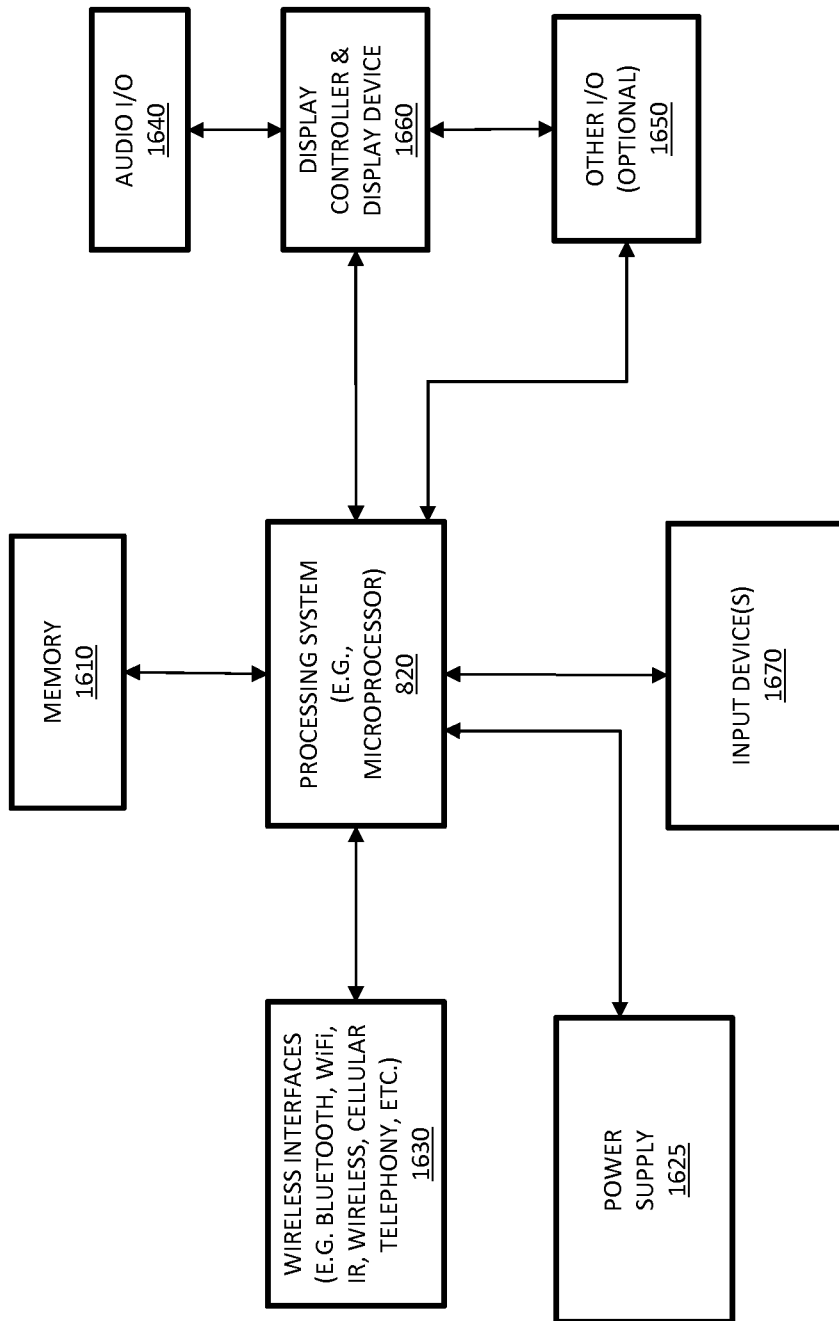


FIG. 14

**FIG. 15**

**FIG. 16**

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# SYSTEM AND METHOD FOR IMPLEMENTING PRIVACY CLASSES WITHIN AN AUTHENTICATION FRAMEWORK

## BACKGROUND

### 1. Field of the Invention

This invention relates generally to the field of data processing systems. More particularly, the invention relates to an improved system and method for implementing privacy classes within an authentication framework.

### 2. Description of Related Art

Existing systems have been designed for providing secure user authentication over a network using biometric sensors. In particular, an Online Secure Transaction Plugin (OSTP) protocol developed by the Fast Identify Online (FIDO) alliance enables strong authentication (e.g., protection against identity theft and phishing), secure transactions (e.g., protection against “malware in the browser” and “man in the middle” attacks for transactions), and enrollment/management of client authentication tokens (e.g., fingerprint readers, facial recognition devices, smartcards, trusted platform modules, etc). Details of the existing OSTP protocol can be found, for example, in U.S. Patent Application No. 2011/0082801 (“’801 Application”), and the document entitled OSTP Framework (Mar. 23, 2011), both of which describe a framework for user registration and authentication on a network.

## BRIEF DESCRIPTION OF THE DRAWINGS

A better understanding of the present invention can be obtained from the following detailed description in conjunction with the following drawings, in which:

FIGS. 1A-B illustrate two different embodiments of a secure authentication system architecture.

FIG. 2 is a transaction diagram showing how authentication devices on a client device may be discovered.

FIG. 3 is a transaction diagram showing how a user may enroll with authentication devices.

FIG. 4 is a transaction diagram showing how keys may be registered into authentication devices.

FIG. 5 is a transaction diagram showing how user authentication may be implemented within an authentication framework.

FIG. 6 is a transaction diagram showing how details of a transaction may be verified.

FIG. 7 illustrates a query policy filter implemented in accordance with one embodiment of the invention.

FIG. 8 is a transaction diagram showing how a registration operation with query policy is implemented in one embodiment of the invention.

FIG. 9 illustrates one embodiment of an architecture for implementing multiple authentication device processing.

FIG. 10A-C illustrate three embodiments of the invention for multiple authentication device processing.

FIGS. 11A-B illustrate a transaction diagram for detecting and responding to a random challenge timeout.

FIG. 12 illustrates an architecture for implementing privacy classes in accordance with one embodiment of the invention.

FIG. 13 is a transaction diagram for implementing privacy classes in accordance with one embodiment of the invention.

FIG. 14 illustrates one embodiment of an architecture for using signatures to authenticate and a transaction.

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FIGS. 15-16 illustrate exemplary embodiments of a computer system for executing embodiments of the invention.

## DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

Described below are embodiments of an apparatus, method, and machine-readable medium for intelligently implementing an authentication framework in a client-server environment. Throughout the description, for the purposes of explanation, numerous specific details are set forth in order to provide a thorough understanding of the present invention. It will be apparent, however, to one skilled in the art that the present invention may be practiced without some of these specific details. In other instances, well-known structures and devices are not shown or are shown in a block diagram form to avoid obscuring the underlying principles of the present invention.

The embodiments of the invention discussed below involve client devices with authentication capabilities such as biometric devices. These devices are sometimes referred to herein as “tokens.” Various different biometric devices may be used including, but not limited to, fingerprint sensors, voice recognition hardware/software (e.g., a microphone and associated software for recognizing a user’s voice), facial recognition hardware/software (e.g., a camera and associated software for recognizing a user’s face), and optical recognition capabilities (e.g., an optical scanner and associated software for scanning the retina of a user). The authentication capabilities may also include non-biometric devices such as a trusted platform modules (TPMs) and smartcards.

The embodiments of the invention described below provide various improvements over existing authentication techniques. For example, in contrast to current techniques which require a client to communicate an exhaustive list of all of its authentication capabilities (e.g., all of its authentication tokens/devices) over a network, one embodiment of the invention implements a query policy in which a secure transaction server initially transmits a server policy to the client indicating the authentication capabilities accepted by the server. The client then analyzes the server policy to identify a subset of authentication capabilities, thereby reducing the privacy impact to the client.

In another embodiment, multiple configurable levels of privacy protection are employed. Privacy classes are pre-defined and may be selected and/or modified by the end user. In one embodiment of the invention, the privacy classes are defined based on the probability with which a client can be identified using the information requested. At relatively higher privacy levels (having a relatively lower privacy impact), relatively less information about the client device is divulged to perform the authentication techniques described herein.

Another embodiment of the invention provides for the provisioning or authentication of multiple devices at the same time, thereby improving efficiency. For example, instead of requesting registration or authentication for a single authentication device at a time, a list of authentication devices may be sent from the server. Symmetric and/or asymmetric keys are then provisioned into multiple tokens/devices in one operation, or series of sequential operations executed locally on the client. For authentication, several tokens/devices may be selected concurrently for a given transaction.

Another embodiment of the invention improves the efficiency with which server challenges are processed and managed. Today, after a server sends a random challenge to the client (e.g., a cryptographic nonce), if the client does not

respond within a specified timeout period, the nonce is no longer valid and the client will receive an error in response to a subsequent authentication attempt. For example, if the user suspends the client to move to a new location (e.g., closing the lid on a laptop) and then attempts authentication, the authentication attempt will be denied. In one embodiment of the invention, the client detects that the random challenge has expired and automatically and transparently requests a new challenge from the server. The server then generates a new challenge and transmits it to the client where it may be used for authentication. The end user experience is improved because the user does not receive an error or denial of an authentication request.

Another embodiment of the invention employs transaction signing on a secure transaction server so that no transaction state needs to be maintained on the server to maintain current sessions with clients. Transaction content such as transaction text is sent to the client signed by server, when the server responds it sends back the transaction content with the signature. The server does not need to store transaction state because it can verify that the signed transaction responses received by the client are valid by verifying the signature.

While described above as separate embodiments, all of the above techniques may be combined together in various ways within a single comprehensive authentication system. Thus, a given embodiment of the invention may be combined with one or more other embodiments described herein for improving client and user authentication in a secure network environment.

#### Exemplary System Architectures

FIGS. 1A-B illustrate two embodiments of a system architecture comprising client-side and server-side components for authenticating a user. The embodiment shown in FIG. 1A uses a browser plugin-based architecture for communicating with a website while the embodiment shown in FIG. 1B does not require a browser. The various techniques described herein such as enrolling a user with authentication devices, registering the authentication devices with a secure server, and authenticating a user may be implemented on either of these system architectures. Thus, while the architecture shown in FIG. 1A is used to demonstrate the operation of several of the embodiments described below, the same basic principles may be easily implemented on the system shown in FIG. 1B (e.g., by removing the browser plugin 105 as the intermediary for communication between the server 130 and the secure transaction service 101 on the client).

Turning first to FIG. 1A, the illustrated embodiment includes a client 100 equipped with one or more authentication devices 110-112 (sometimes referred to in the art as authentication “tokens”) for enrolling and authenticating an end user. As mentioned above, the authentication devices 110-112 may include biometric device such as fingerprint sensors, voice recognition hardware/software (e.g., a microphone and associated software for recognizing a user’s voice), facial recognition hardware/software (e.g., a camera and associated software for recognizing a user’s face), and optical recognition capabilities (e.g., an optical scanner and associated software for scanning the retina of a user) and non-biometric devices such as a trusted platform modules (TPMs) and smartcards.

The authentication devices 110-112 are communicatively coupled to the client through an interface 102 (e.g., an application programming interface or API) exposed by a secure transaction service 101. The secure transaction service 101 is a secure application for communicating with one or more

secure transaction servers 132-133 over a network and for interfacing with a secure transaction plugin 105 executed within the context of a web browser 104. As illustrated, the Interface 102 may also provide secure access to a secure storage device 120 on the client 100 which stores information related to each of the authentication devices 110-112 such as a device identification code, user identification code, user enrollment data (e.g., scanned fingerprint or other biometric data), and keys used to perform the secure authentication techniques described herein. For example, as discussed in detail below, a unique key may be stored into each of the authentication devices and used when communicating to servers 130 over a network such as the Internet.

As discussed below, certain types of network transactions are supported by the secure transaction plugin 105 such as HTTP or HTTPS transactions with websites 131 or other servers. In one embodiment, the secure transaction plugin is initiated in response to specific HTML tags inserted into the HTML code of a web page by the web server 131 within the secure enterprise or Web destination 130 (sometimes simply referred to below as “server 130”). In response to detecting such a tag, the secure transaction plugin 105 may forward transactions to the secure transaction service 101 for processing. In addition, for certain types of transactions (e.g., such as secure key exchange) the secure transaction service 101 may open a direct communication channel with the on-premises transaction server 132 (i.e., co-located with the website) or with an off-premises transaction server 133.

The secure transaction servers 132-133 are coupled to a secure transaction database 120 for storing user data, authentication device data, keys and other secure information needed to support the secure authentication transactions described below. It should be noted, however, that the underlying principles of the invention do not require the separation of logical components within the secure enterprise or web destination 130 shown in FIG. 1A. For example, the website 131 and the secure transaction servers 132-133 may be implemented within a single physical server or separate physical servers. Moreover, the website 131 and transaction servers 132-133 may be implemented within an integrated software module executed on one or more servers for performing the functions described below.

As mentioned above, the underlying principles of the invention are not limited to a browser-based architecture shown in FIG. 1A. FIG. 1B illustrates an alternate implementation in which a stand-alone application 154 utilizes the functionality provided by the secure transaction service 101 to authenticate a user over a network. In one embodiment, the application 154 is designed to establish communication sessions with one or more network services 151 which rely on the secure transaction servers 132-133 for performing the user/client authentication techniques described in detail below.

In either of the embodiments shown in FIGS. 1A-B, the secure transaction servers 132-133 may generate the keys which are then securely transmitted to the secure transaction service 101 and stored into the authentication devices within the secure storage 120. Additionally, the secure transaction servers 132-133 manage the secure transaction database 120 on the server side.

#### An Overview of Device Discovery, Enrollment, Registration, and Authentication

An exemplary series of transactions for performing authentication device discovery, enrollment, registration, and authentication are shown in FIGS. 2-6. Some aspects of these



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transactions have been employed in the OSTP protocol mentioned above (see the OSTP Framework (Mar. 23, 2011) for additional details, which is incorporated herein by reference). An understanding of the basic operation of these transactions will provide a context in which embodiments of the invention may be implemented.

The operations described below include detection of authentication devices (FIG. 2); enrollment of the user with the authentication devices (FIG. 3); registration of authentication devices (FIG. 4); user authentication with the registered authentication devices (FIG. 5); and implementation of secure transactions following authentication (FIG. 6).

FIG. 2 illustrates a series of transactions for detecting authentication devices on the client machine. After device detection is successfully completed, the server 130 possesses exhaustive information about the authentication devices attached to the client and will be able to assess which device(s) are most appropriate to use with the enhanced security infrastructure. Only the server 130 filters the list of authentication devices. The user will be provided with this list and may choose one (or combination) of authentication devices to use for further authentication and implementation of secure transactions.

In operation, the user authenticates with username and password in browser and logs in to web site. This is the only time that the user will be required to provide a user name and password. The server 130 determines that the user is not currently using enhanced security (e.g., by querying the secure transaction database 120) and provides a suggestion to the user to change to enhanced security.

In one embodiment, the server 130 includes a “query for devices” tag in an HTML page which the secure transaction plugin 105 detects. In response to detecting the tag, the secure transaction plugin 105 reroutes the request to the secure transaction service 101 which then prepares exhaustive information about all authentication devices attached to the system including security characteristics of the devices. In one embodiment, the information is packaged in an XML format prior to transmission using a pre-specified data schema.

The secure transaction plugin 105 receives this information from the secure transaction service 101 and, in one embodiment, passes the information to the web page’s JavaScript via a registered callback. It then chooses how to display the information in the browser 104. The list, filtered by the website, may be shown to the user and the user may select one or a combination of authentication devices.

FIG. 3 illustrates a series of transactions to enroll the user with the authentication devices. In one embodiment, enrollment is a prerequisite for using the enhanced security provided by the embodiments of the invention described herein. Enrollment involves taking a biometric reading of the user (e.g., a fingerprint, voice sample, etc) so that the same authentication device can be used to authenticate the user during a subsequent transaction. The enrollment operation may be done solely on the client, without interaction with the server 130. The user interface(s) provided for enrollment may be displayed in the browser extension or may be displayed in a separate application or mobile device app.

The enrollment operation may be initiated as soon as devices are detected. The user may choose to use one or a group of discovered devices for enhanced security. In operation, the user may select a device from the displayed device list in the browser, application or mobile device app. For the browser-based implementation illustrated in FIG. 3, the secure transaction plugin 105 displays a device-specific enrollment graphical user interface (GUI). The secure transaction plugin 105 transmits the device identifier and an enroll-

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ment request to secure transaction service 101 and waits for completion. If the user is already enrolled with an authentication device on the client, the user may only need to verify their identity (i.e., they will not be required to enroll again). If the user is not currently enrolled, then the secure transaction service 101 starts the enrollment process by activating the physical authentication device (e.g., via the device interface 102). The user then interacts with the secure transaction plugin 105 GUI and follows the specified enrollment steps (e.g., swiping a finger, speaking into a microphone, snapping a picture, etc). Once complete, the user will be enrolled with the authentication device. Significantly, once a user is enrolled with a device, they may use this enrollment to register or authenticate with any website or network service as described herein.

FIG. 4 illustrates a series of transactions for registration of authentication devices. During registration, a key is shared between the authentication device and one of the secure transaction servers 132-133. The key is stored within the secure storage 120 of the client 100 and the secure transaction database 120 used by the secure transaction servers 132-133. In one embodiment, the key is a symmetric key generated by one of the secure transaction servers 132-133. However, in another embodiment discussed below, asymmetric keys may be used. In this embodiment, the public key may be stored by the secure transaction servers 132-133 and a second, related private key may be stored in the secure storage 120 on the client. Moreover, in one embodiment (also discussed below), the key(s) may be generated on the client 100 (e.g., by the authentication device or the authentication device interface rather than the secure transaction servers 132-133).

A secure key provisioning protocol such as the Dynamic Symmetric Key Provisioning Protocol (DSKPP) may be used to share the key with the client over a secure communication channel (see, e.g., Request for Comments (RFC) 6063). However, the underlying principles of the invention are not limited to any particular key provisioning protocol.

Turning to the specific details shown in FIG. 4, once the user enrollment or user authentication is complete, the server 130 generates a randomly generated challenge (e.g., a cryptographic nonce) that must be presented by the client during device registration. The random challenge may be valid for a limited period of time. The secure transaction plugin detects the random challenge and forwards it to the secure transaction service 101. In response, the secure transaction service initiates an out-of-band session with the server 130 (e.g., an out-of-band transaction) and communicates with the server 130 using the key provisioning protocol. The server 130 locates the user with the user name, validates the random challenge, validates the device’s authentication code if one was sent, and creates a new entry in the secure transaction database 120 for the user. It may also generate the key, write the key to the database 120 and send the key back to the secure transaction service 101 using the key provisioning protocol. Once complete, the authentication device and the server 130 share the same key if a symmetric key was used or different keys if asymmetric keys were used.

FIG. 5 illustrates a series of transactions for user authentication with the registered authentication devices. Once device registration is complete the server 130 will accept a token generated by the local authentication device as a valid authentication token.

Turning to the specific details shown in FIG. 5, which shows a browser-based implementation, the user enters the uniform resource locator (URL) of the server 130 in the browser 104. In an implementation which uses a stand alone application or mobile device app (rather than a browser), the

user may enter a network address for a network service or the application or app may automatically attempt to connect to the network service at the network address.

For a browser-based implementation, the website embeds a query for registered devices in the HTML page. This may be done in many ways other than embedding the query in an HTML page, such as through Javascript or using HTTP headers. The secure transaction plugin **105** receives the URL and sends it to secure transaction service **101**, which searches the looks into the secure storage **120** (which, as discussed, includes a database of authentication device and user information) and determines whether there is a user enrolled within this URL. If so, the secure transaction service **101** sends a list of provisioned devices associated with this URL to the secure transaction plugin **105**. The secure transaction plugin then calls the registered JavaScript API and passes this information to the server **130** (e.g., the website). The server **130** chooses the appropriate device from the sent device list, generates a random challenge and sends the device information, and argument back to the client. The website displays the corresponding user interface and asks for authentication from the user. The user then provides the requested authentication measure (e.g., swiping a finger across the fingerprint reader, speaking for voice recognition, etc). The secure transaction service **101** identifies the user (this step can be skipped for devices which don't support storing users), obtains the username from the database, generates an authentication token using the key and sends this information to the website via the secure transaction plugin. The server **130** identifies the user from the secure transaction database **120** and verifies the token by generating the same token on the server **130** (e.g., using its copy of the key). Once verified, the authentication process is complete.

FIG. 6 illustrates a secure transaction following authentication for a browser-based implementation. The secure transaction is designed to provide stronger security for certain types of transactions (e.g., financial transactions). In the illustrated embodiment, the user confirms each transaction prior to committing the transaction. Using the illustrated techniques, the user confirms exactly what he/she wants to commit and commits exactly what he/she sees displayed in the GUI. In other words, this embodiment ensures that the transaction text cannot be modified by a "man in the middle" to commit a transaction which the user did not confirm.

In one embodiment, the secure transaction plugin **105** displays a window **601** in the browser context to show the transaction details. The secure transaction server **101** periodically (e.g., with a random interval) verifies that the text that is shown in the window is not being tampered by anyone.

The following example will help to highlight the operation of this embodiment. A user chooses items for purchase from a merchant site and selects "check out." The merchant site sends the transaction to a service provide which has a secure transaction server **132-133** implementing one or more of the embodiments of the invention described herein (e.g., PayPal). The merchant site authenticates the user and completes the transaction.

The secure transaction server **132-133** receives the transaction details (TD) and puts a "Secure Transaction" request in an HTML page and sends to client **100**. The Secure Transaction request includes the transaction details and a random challenge (e.g., a random nonce). The secure transaction plugin **105** detects the request for transaction confirmation message and forwards all data to the secure transaction service **101**. In an embodiment which does not use a browser or

plugin, the information may be sent directly from the secure transaction servers to the secure transaction service on the client **100**.

For a browser-based implementation, the secure transaction plugin **105** displays a window **601** with transaction details to the user (in a browser context) and asks the user to provide authentication to confirm the transaction. In an embodiment which does not use a browser or plugin, the secure transaction service **101** or application **154** may display the window **601**. The secure transaction service **101** starts a timer and verifies the content of the window **601** being displayed to the user. The period of verification may be randomly chosen. The secure transaction service **101** ensures that user sees the valid transaction details in the window **601**. If it detects that the content has been tampered with it prevents the confirmation token from being generated.

After the user provides valid authentication (e.g., swipes a finger on the fingerprint sensor), the device identifies the user and generates a token (cryptographic signature) with the transaction details and the random challenge (i.e., the token is calculated over the transaction details and the nonce). This allows the secure transaction server **132-133** to ensure that the transaction details have not been modified between the server and the client. The secure transaction service **101** sends the generated token and username to the secure transaction plugin **105** which forwards the token to the secure transaction server **132-133**. The secure transaction server **132-133** identifies the user with the username and verifies the token. If verification succeeds, a confirmation message is sent to the client and the transaction is processed.

#### System and Method for a Secure Query Policy to Determine Client Authentication Capabilities

As mentioned, one embodiment of the invention implements a query policy in which a secure transaction server transmits a server policy to the client indicating the authentication capabilities accepted by the server. The client then analyzes the server policy to identify a subset of authentication capabilities which it supports and/or which the user has indicated a desire to use. The client then registers and/or authenticates the user using the subset of authentication tokens matching the provided policy. Consequently, there is a lower impact to the client's privacy because the client is not required to transmit exhaustive information about its authentication capabilities (e.g., all of its authentication devices) or other information which might be used to uniquely identify the client.

By way of example, and not limitation, the client may include numerous authentication capabilities such as a fingerprint sensor, voice recognition capabilities, facial recognition capabilities, eye/optical recognition capabilities, a trusted platform module (TPM), and smartcard, to name a few. However, for privacy reasons, the user may not wish to divulge the details for all of its capabilities to a requesting server. Thus, using the techniques described herein, the secure transaction server may transmit a server policy to the client indicating that it supports, for example, fingerprint, optical, or smartcard authentication. The client may then compare the server policy against its own authentication capabilities and choose one or more of the available authentication options.

FIG. 7 illustrates one embodiment of a client-server architecture for implementing these techniques. As illustrated, the secure transaction service **101** implemented on the client **100** includes a policy filter **701** for analyzing the policy provided by the server **130** and identifying a subset of authentication

capabilities to be used for registration and/or authentication. In one embodiment, the policy filter **701** is implemented as a software module executed within the context of the secure transaction service **101**. It should be noted, however, that the policy filter **701** may be implemented in any manner while still complying with the underlying principles of the invention and may include software, hardware, firmware, or any combination thereof.

The particular implementation shown in FIG. 7 includes a secure transaction plugin **105** for establishing communication with the secure enterprise or Web destination **130** (sometimes referred to simply as “server **130**”) using techniques previously discussed. For example, the secure transaction plugin may identify a specific HTML tag inserted into the HTML code by a web server **131**. Thus, in this embodiment, the server policy is provided to the secure transaction plugin **105** which forwards it to the secure transaction service **101** implementing the policy filter **701**.

The policy filter **701** may determine the client authentication capabilities by reading the capabilities from the client’s secure storage area **720**. As previously discussed, the secure storage **720** may comprise a repository of all of the client’s authentication capabilities (e.g., identification codes for all of the authentication devices). If the user has already enrolled the user with its authentication devices, the user’s enrollment data is stored within the secure storage **720**. If the client has already registered an authentication device with a server **130**, then the secure storage may also store an encrypted secret key associated with each authentication device.

Using the authentication data extracted from the secure storage **720** and the policy provided by the server, the policy filter **701** may then identify a subset of authentication capabilities to be used. Depending on the configuration, the policy filter **701** may identify a complete list of authentication capabilities supported by both the client and the server or may identify a subset of the complete list. For example, if the server supports authentication capabilities A, B, C, D, and E and the client has authentication capabilities A, B, C, F, and G, then the policy filter **701** may identify the entire subset of common authentication capabilities to the server: A, B, and C. Alternatively, if a higher level of privacy is desired, as indicated by user preferences **730** in FIG. 7, then a more limited subset of authentication capabilities may be identified to the server. For example, the user may indicate that only a single common authentication capability should be identified to the server (e.g., one of A, B or C). In one embodiment, the user may establish a prioritization scheme for all of the authentication capabilities of the client **100** and the policy filter may select the highest priority authentication capability (or a prioritized set of N authentication capabilities) common to both the server and the client.

Depending on what operation has been initiated by server **130** (Registration or Authentication), the secure transaction service **130** performs that operation on the filtered subset of authentication devices (**110-112**) and sends the operation response back to server **130** via the secure transaction plugin **105** as shown in FIG. 7. Alternatively, in an embodiment which does not rely on a plugin **105** component of a Web browser, the information may be passed directly from the secure transaction service **101** to the server **130**.

FIG. 8 illustrates a transaction diagram showing additional details for an exemplary series of registration with query policy transaction. In the illustrated embodiment, the user has not previously registered devices with the server **130**. Consequently, at **801** the user may enter a user name and password as an initial, one-time authentication step, which is forwarded at **802** to the server **130** via the client browser **104**. It should

be noted, however, that a user name and password are not required for complying with the underlying principles of the invention.

Because the user has not previously registered with enhanced security, determined at **803**, the server **130** transmits its server policy to the client at **804**. As mentioned, the server policy may include an indication of the authentication capabilities supported by the server **130**. In the illustrated example, the server policy is passed to the secure transaction service **101** via transaction **806**.

At transaction **807**, the secure transaction service **101** compares the server policy with the capabilities of the client (and potentially other information such as device priority scheme and/or user preferences as described above) to arrive at a filtered list of authentication capabilities. The filtered list of devices (**102**) then generate keys (**808** and **809**) and then provide public parts of these keys to secure transaction service **101** which in its turn sends these as registration response back to server **130**. The server attests the authentication devices and stores public keys in secure transaction database. The Token Attestation employed here is the process of validating authentication device identity during registration. It allows server to cryptographically make sure that the device reported by Client is really who it claimed to be.

Alternatively, or in addition, at **807**, the user may be provided with an opportunity to review the list and/or select specific authentication capabilities to be used with this particular server **130**. For example, the filtered list may indicate the option to use authentication with a fingerprint scan, facial recognition, and/or voice recognition. The user may then choose to use one or more of these options when authenticating with the server **130**.

The techniques described above for filtering a server policy at a client may be implemented at various different stages of the series of transactions described above (e.g., during device discovery, device registration, device provisioning, user authentication, etc). That is, the underlying principles of the invention are not limited to the specific set of transactions and the specific transaction ordering set forth in FIG. 8.

Moreover, as previously mentioned, a browser plugin architecture is not required for complying with the underlying principles of the invention. For an architecture which does involve a browser or browser plug-ins (e.g., such as a stand-alone application or mobile device app), the transaction diagram shown in FIG. 8 (and the rest of the transaction diagrams disclosed herein) may be simplified such that the browser **104** is removed, and the secure transaction service **101** communicates directly with the server **130**.

#### System and Method for Efficiently Enrolling, Registering, and Authenticating with Multiple Authentication Devices

One embodiment of the invention is capable of enrolling, registering, and authenticating multiple devices at the same time, thereby improving efficiency and the user experience. For example, instead of requesting registration and authentication for a single device at a time, a list of devices may be sent to the client. Symmetric or asymmetric keys may then be registered into multiple devices in one operation, or series of sequential operations executed locally on the client. For authentication, several tokens/devices may be selected concurrently for a given transaction.

FIG. 9 illustrates one embodiment of a client-server architecture for implementing these techniques. As illustrated, the secure transaction service **101** implemented on the client **100** includes multi-device processing logic **901** for performing

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specified operations such as enrollment and registration of multiple devices at a time without the need for continual back-and-forth communication with the server 130 as each device is enrolled/registered. Similarly, the server 130 includes multi-device processing logic for issuing commands directed to multiple authentication devices. In one embodiment, the multi-device processing logic 901 is implemented as a software module executed within the context of the secure transaction service 101. It should be noted, however, that the multi-device processing logic 901 may be implemented in any manner while still complying with the underlying principles of the invention and may include software, hardware, or firmware components, or any combination thereof.

As in the embodiments described above, the particular implementation shown in FIG. 9 includes secure transaction plugin 105 for establishing communication with the server 130 (which, as discussed, may include a Website server 131 and secure transaction servers 132-133). Thus, the server 130 communicates with the secure transaction service 101 via the secure transaction plugin 105. As mentioned, however, a browser-based plugin architecture is not required for complying with the underlying principles of the invention.

The multi-device processing logic 902 on the server 130 may communicate commands to be executed by the multi-device processing logic 901 on the client 100 which performs the operations on multiple authentication devices 110-112. By way of example, the multi-device processing logic 902 may generate N keys to be registered with each of N authentication devices and then transmit securely to the multi-device processing logic 901 along with a command to register the N devices. The multi-device processing logic 901 may then perform the registration concurrently or in a series of sequential operations for all N devices (e.g., for authentication devices 110-112) without further interaction with the server. A single response may then be sent to the server 130 to indicate the completed registration of all N devices.

A series of exemplary multiple transactions are illustrated in FIGS. 10A-C. FIG. 10A illustrates a multiple-device enrollment process which may be performed without any interaction with the server 130 (e.g., enrolling the user with authentication devices may be performed under the control of the secure transaction service 101 on the client). In an alternate embodiment, the server 130 may transmit a request to the client (now shown) to enroll the user with the N devices. FIGS. 10B-C illustrate two different embodiments for registering multiple devices with the server 130.

Turning to the enrollment process in FIG. 10A, at 1001, the user indicates a desire to enroll with N authentication devices on the client (representing all or a subset of the available authentication devices). In response, the secure transaction plugin is called at 1002 and, at 1003, a device-specific graphical user interface (GUI) is generated to walk the user through the process of enrolling with authentication device #1. During the enrollment process, the user interacts with the secure transaction plugin as indicated (e.g., by positioning a finger over a fingerprint sensor, speaking into a microphone, snapping a picture with the camera, etc). In one embodiment, enrollment is performed for each of the N devices until enrollment is completed for the Nth device at 1004. A different, device-specific script and/or user interface may be presented to the user to enroll the user with each individual authentication device. As previously discussed, as the user enrolls with each device, the user enrollment data may be stored within the secure storage 720 on the client 100 and made accessible only through the secure transaction service 101. Once enrollment

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for all N devices is complete, a notification may be sent to the server 130 via transactions 1004-1005.

Regardless of how enrollment is performed, once completed, the transaction diagram shown in FIG. 10B may be used to register the N devices with the server 130. At 1010 the server 130 generates a user-specific random challenge which, as previously described, may only be valid for a limited window of time and may comprise a randomly-generated code such as a cryptographic nonce. At 1011, the random challenge is transmitted along with a command to register N authentication devices with the server 130. At 1012, the secure transaction service 101 creates a secure connection with the server 130 and transmits identification data for the N devices, along with the random challenge. In one embodiment, the secure connection is an HTTPS connection. However, the underlying principles of the invention are not limited to any particular secure connection type.

At 1013, the server 130 attests the N devices, generates a key for each of the N devices, and sends the N keys back to the secure transaction service over the secure connection. In one embodiment, the Dynamic Symmetric Key Provisioning Protocol (DSKPP) is used to exchange keys with the client over the secure connection. However, the underlying principles of the invention are not limited to any particular key provisioning techniques. Alternatively, in an embodiment which does not rely on DSKPP protocol, the keys may be generated in each Authentication Device and then transmitted to server 130.

At 1014-1015, the multi-device processing logic of the secure transaction service registers each of the N keys into each of the N devices. As previously described, each key may be stored and associated with its respective device within the secure storage 720 on the client. Once registration is complete for each authentication device, a notification is sent to the server over the secure connection at 1016.

In one embodiment, the keys registered into each authentication device are symmetric keys. Thus, an identical copy of each key is stored in the secure storage 720 on the client and the secure transaction database 120 on the server 130. In an alternate implementation, asymmetric key pairs may be generated, with one of the keys being maintained as a public key in the secure transaction database 120 on the server and the private key being stored in the secure storage 720 of the client. It should be noted, however, that the underlying principles of the invention are not limited to any particular type of encryption keys.

An alternate implementation is illustrated in FIG. 10C in which keys are generated on the client rather than the server 130. In this implementation, after receiving the request to register devices with the random challenge at 1011, the multi-device processing logic of the secure transaction service 101 generates N keys for each of the N devices at 1120. Once generated, the keys are registered with each of the N devices at 1013-1014 and the registration stored within the secure storage 720 as previously described. Once all keys have been registered, the secure transaction service 101 provides a notification to the server at 1015 along with the random challenge (to verify the identity of the client). The server 130 may then store the registration in the secure transaction database 120 as described above.

#### System and Method for Processing Random Challenges within an Authentication Framework

One embodiment of the invention improves the manner in which random challenges are generated by the server and processed. In one embodiment, the random challenge com-

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prises a randomly generated code such as a cryptographic nonce. In current systems, after a server transmits a random challenge to the client, if the client does not respond within a specified timeout period, the random challenge is no longer valid and the client will receive an error in response to a subsequent authentication attempt (e.g., the user will swipe a finger on the fingerprint reader and be denied).

In one embodiment of the invention, the client automatically detects that the challenge has expired and transparently requests a new challenge from the server (i.e., without user intervention). The server then generates a new random challenge and transmits it to the client which may then use it to establish secure communication with the server. The end user experience is improved because the user does not receive an error or denial of an authentication request.

FIG. 11A illustrates one such embodiment which is used within the context of a registration process and FIG. 11B illustrates an embodiment which is used within the context of an authentication process. It should be noted, however, that the underlying principles of the invention may be employed in other contexts than those shown in FIGS. 11A-B. For example, the techniques described herein may be used with any process in which a time-sensitive code is communicated from a server to a client.

Turning first to FIG. 11A, at 1101, the server 130 generates a random challenge and an indication of a timeout period. In one embodiment, the timeout period comprises a period of time for which the random challenge is considered valid. After the timeout period has elapsed, the random challenge is no longer considered valid by the server 130. In one embodiment, the timeout period is specified simply as a point in time at which the random challenge will no longer be valid. Once this point in time is reached, the random challenge is invalid. In another embodiment, the timeout period is specified by using a current timestamp (i.e., the time at which the random challenge is generated by the server 130) and a duration. The secure transaction service 101 may then calculate the timeout time by adding the duration value to the timestamp to calculate the point in time when the random challenge becomes invalid. It should be noted, however, that the underlying principles of the invention are not limited to any specific technique for calculating the timeout period.

Regardless of how the timeout period is specified or calculated, at 1102 the random challenge and the timeout indication are transmitted to the secure transaction service 101 (via the browser 104 and secure transaction plugin 105 in the illustrated example). At 1103, the secure transaction service 101 detects that the random challenge has timed out and is no longer valid based on the timeout indication sent from the server 130. By way of example, the user may have turned off his/her client machine or closed the lid on his/her notebook computer prior to completing the series of transactions. If the transaction is one which requires user interaction, the user may have simply walked away or ignored a message displayed within the GUI.

At 1104, upon detecting that the random challenge is no longer valid, the secure transaction service 101 transmits a request for a new random challenge to the server 130 (via the secure transaction plugin 105 and browser 104 in the illustrated example). At 1105, the server 130 generates a new random challenge and a new indication of the timeout period. In one embodiment, the timeout period is the same as in operation 1101 or may be modified. For example, the server 130 may increase the duration of the timeout period to reduce data traffic with the client or decrease the duration to increase the level of security provided by the random challenge. At

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1106, the new random challenge and timeout indication is transmitted to the secure transaction service 101.

The remainder of the transactions occurs as previously described. For example, the secure transaction service opens a secure connection directly to the server at 1107 in order to perform device registration and key exchange as discussed above with respect to FIG. 4, FIG. 10B, or FIG. 10C. At 1108, the server 130 identifies the user (e.g., with a user name or other ID), attests authentication device, and generates a key for the device. As mentioned, the key may be a symmetric key or an asymmetric key. At 1109, the keys are transmitted to the secure transaction service 101 via the secure connection and, at 1110, the secure transaction service 101 registers the key into the authentication device. At 1111, a notification that registration is complete is transmitted to the server 130.

Thus, in the embodiment shown in FIG. 11A, the key used for device registration is generated at the server 130 as in the embodiment shown in FIG. 10B. However, the underlying principles of the invention may also be used in an embodiment in which the key(s) are generated by the secure transaction service 101 on the client 100, such as that described above with respect to FIG. 10C.

FIG. 11B illustrates one embodiment of the invention implemented within the context of an authentication process. At 1151, the user enters a particular website URL into the browser 104 and is directed to the web server 131 within the enterprise/web destination servers 130 which includes the secure transaction servers 132-133. At 1152, a query is sent back to the secure transaction service (via the browser and plugin) to determine which device(s) are registered with the website's URL. The secure transaction service 101 queries the secure storage 720 on the client 100 to identify a list of devices which are sent back to the server 130 at 1153. At 1154, the server 1154 chooses a device to use for authentication, generates a random challenge and a timeout indication and, at 1155, sends this information back to the secure transaction service 101.

At 1156, the secure transaction service 1156 automatically detects that the random challenge is no longer valid upon reaching the end of the timeout period. As mentioned above, various different techniques may be employed for indicating and detecting the end of the timeout period (see FIG. 11A and associated text). Upon detecting the expiration of the random challenge, at 1157, the secure transaction service 101 transparently (i.e., without user intervention) notifies the server 130 and requests a new random challenge. In response, at 1158, the server 130 generates a new random challenge and a new indication of the timeout period. As mentioned, the new timeout period may be the same as previously sent to the client or may be modified. In either case, at 1159, the new random challenge and timeout indication are sent to the secure transaction service 101.

The remainder of the transaction diagram shown in FIG. 11B operates in substantially the same manner as described above (see, e.g., FIG. 5). For example, at 1160, an authentication user interface is displayed (e.g., directing the user to swipe a finger on a fingerprint sensor) and, at 1161, the user provides authentication (e.g., swipes a finger on the fingerprint scanner). At 1162, the secure transaction service verifies the identity of the user (e.g., comparing the authentication data collected from the user with that stored in the secure storage 720) and uses the key associated with the authentication device to encrypt the random challenge. At 1163, the user name (or other ID code) and the encrypted random challenge are sent to the server 130. Finally, at 1164, the server 130 identifies the user within the secure transaction database 120 using the user name (or other ID code), and decrypts/verifies

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the random challenge using the key stored in the secure transaction database **120** to complete the authentication process.

#### System and Method for Implementing Privacy Classes within an Authentication Framework

In one embodiment, multiple classes of privacy protection may be predefined, selected and/or modified by the end user. The privacy classes may be defined based on the probability with which a client can be identified using the divulged information. At privacy classes having relatively higher privacy levels, relatively less information about the client device is divulged to perform the authentication techniques described herein. In one embodiment, the user may choose to disclose the least amount of information possible when communicating with different servers (i.e., may choose transactions having the lowest allowable privacy impact for each website or network service).

FIG. **12** illustrates a high level architecture for implementing privacy classes. As illustrated, the secure transaction service **101** of this embodiment includes privacy management logic **1201** for analyzing queries received from the server **130** for client information such as information related to authentication devices, implementing a privacy policy in response to such queries, and generating a response containing client information collected based on the particular privacy class in use. In one embodiment, the privacy management module **1201** is implemented as a software module executed within the context of the secure transaction service **101**. It should be noted, however, that the privacy management module **1201** may be implemented in any manner while still complying with the underlying principles of the invention and may include software, hardware, firmware, or any combination thereof.

The privacy classes utilized by the privacy management logic **1201** may be pre-specified and stored on the client **100** (e.g., within stored within secure storage **720**). In one embodiment, three privacy classes are defined: high privacy impact, medium privacy impact, and low privacy impact. Each privacy class may be defined based on a probability with which the divulged information could be used to uniquely identify a user/client. For example, the information divulged for a low privacy impact transaction may result in a 10% probability of the user or machine being uniquely identified over internet; a medium privacy impact transaction may result in a 50% probability of the user or machine being uniquely identified; and a high privacy impact transaction may result in a 100% probability of the user or machine being uniquely identified. Various other privacy class levels may be defined while still complying with the underlying principles of the invention.

In one embodiment, each relying party (e.g., each website **131** or service **151**) may specify a required privacy class or other privacy threshold. For example, websites and services requiring a heightened level of security may only allow communication in accordance with the high privacy impact class whereas other websites/services may permit interactions using the medium privacy impact or low privacy impact class. In one embodiment, the query for client information sent from the server **130** includes an attribute specifying which privacy classes of information should be retrieved (i.e. low, medium, high). Thus, the privacy management logic **1201** will store information for the highest approved privacy class for each relying party. In one embodiment, whenever the relying party asks for information belonging to a higher privacy class than the one already approved, the user will be prompted to permanently approve (or reject) this new privacy

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class for this relying party. In response to the user's approval, the privacy management logic may store the new association between the relying party (e.g., identified via a URL) and the new privacy class.

While the user preferences **1230** are applied directly to the privacy management logic in FIG. **12** for simplicity, it should be noted that the user may specify preferences via a browser-based graphical user interface (not shown). In such a case, the user would enter privacy setting via a browser window. The secure transaction plugin **105** would then store the new settings to the privacy management logic **1201**, or to a configuration data file accessible by the privacy management logic **1201**. In short, the underlying principles of the invention are not limited to any particular mechanism for configuring the privacy management logic.

Various types of client data may be specified at the various privacy class levels including, for example, a machine model identifier, client software information, client capabilities, and various levels of information related to each authentication device configured on the client device (e.g., device ID codes, vendor ID codes, device class ID, etc). Different combinations of this information may be gathered to determine the percentages specified above defining the different privacy classes.

FIG. **13** illustrates a series of transactions for providing information to a requesting party using defined privacy classes. At **1301** the server **130** generates a notification containing a query for client device information. At **1302**, the query is sent to the client and ultimately received by the secure transaction service **101**. At **1303**, the privacy management logic of the secure transaction service determines a privacy class for the response and collects the necessary information. As mentioned above, N different privacy class levels may be defined and the secure transaction service **101** may choose the one which complies with the requirements of the requesting party while at the same time divulges as little information as possible regarding the client. At **1304**, the collected information is sent to the server **130** and at **1305**, the server uses the information for one or more subsequent transactions with the client.

#### System and Method for Implementing an Authentication Framework Using Transaction Signing

One embodiment of the invention employs transaction signing on the secure transaction server so that no transaction state needs to be maintained on the server to maintain sessions with clients. In particular, transaction details such as transaction text may be sent to the client signed by server. The server may then verify that the signed transaction responses received by the client are valid by verifying the signature. The server does not need to persistently store the transaction content, which would consume a significant amount of storage space for a large number of clients and would open possibility for denial of service type attacks on server.

One embodiment of the invention is illustrated in FIG. **14** which shows a website or other network service (**1401**) initiating a transaction with a client **100**. For example, the user may have selected items for purchase on the website and may be ready to check out and pay. In the illustrated example, the website or service **1401** hands off the transaction to a secure transaction server **1402** which includes signature processing logic **1403** for generating and verifying signatures (as described herein) and authentication logic for performing client authentication **1404** (e.g., using the authentication techniques previously described).

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In one embodiment, the authentication request sent from the secure transaction server **1402** to the client **100** includes the random challenge such as a cryptographic nonce (as described above), the transaction details (e.g., the specific text presented to complete the transaction), and a signature generated by the signature processing logic **1403** over the random challenge and the transaction details using a private key (known only by the secure transaction server).

Once the above information is received by the client, the user may receive an indication that authentication is required to complete the transaction. In response, the user may, for example, swipe a finger across a fingerprint scanner, snap a picture, speak into a microphone, or perform any other type of authentication permitted for the given transaction. In one embodiment, once the user has successfully authenticated on the client **100**, the client transmits the following back to the server: (1) the random challenge and transaction text (both previously provided to the client by the server), (2) authentication data proving that the user successfully completed authentication, and (3) the signature.

The authentication module **1404** on the secure transaction server **1402** may then confirm that the user has correctly authenticated and the signature processing logic **1403** regenerates the signature over the random challenge and the transaction text using the private key. If the signature matches the one sent by the client, then the server can verify that the transaction text is the same as it was when initially received from the website or service **1401**. Storage and processing resources are conserved because the secure transaction server **1402** is not required to persistently store the transaction text (or other transaction data) within the secure transaction database **120**.

#### Exemplary Data Processing Devices

FIG. **15** is a block diagram illustrating an exemplary clients and servers which may be used in some embodiments of the invention. It should be understood that while FIG. **15** illustrates various components of a computer system, it is not intended to represent any particular architecture or manner of interconnecting the components as such details are not germane to the present invention. It will be appreciated that other computer systems that have fewer components or more components may also be used with the present invention.

As illustrated in FIG. **15**, the computer system **1500**, which is a form of a data processing system, includes the bus(es) **1550** which is coupled with the processing system **1520**, power supply **1525**, memory **1530**, and the nonvolatile memory **1540** (e.g., a hard drive, flash memory, Phase-Change Memory (PCM), etc.). The bus(es) **1550** may be connected to each other through various bridges, controllers, and/or adapters as is well known in the art. The processing system **1520** may retrieve instruction(s) from the memory **1530** and/or the nonvolatile memory **1540**, and execute the instructions to perform operations as described above. The bus **1550** interconnects the above components together and also interconnects those components to the optional dock **1560**, the display controller & display device **1570**, Input/Output devices **1580** (e.g., NIC (Network Interface Card), a cursor control (e.g., mouse, touchscreen, touchpad, etc.), a keyboard, etc.), and the optional wireless transceiver(s) **1590** (e.g., Bluetooth, WiFi, Infrared, etc.).

FIG. **16** is a block diagram illustrating an exemplary data processing system which may be used in some embodiments of the invention. For example, the data processing system **1600** may be a handheld computer, a personal digital assistant (PDA), a mobile telephone, a portable gaming system, a

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portable media player, a tablet or a handheld computing device which may include a mobile telephone, a media player, and/or a gaming system. As another example, the data processing system **1600** may be a network computer or an embedded processing device within another device.

According to one embodiment of the invention, the exemplary architecture of the data processing system **1600** may be used for the mobile devices described above. The data processing system **1600** includes the processing system **1620**, which may include one or more microprocessors and/or a system on an integrated circuit. The processing system **1620** is coupled with a memory **1610**, a power supply **1625** (which includes one or more batteries) an audio input/output **1640**, a display controller and display device **1660**, optional input/output **1650**, input device(s) **1670**, and wireless transceiver(s) **1630**. It will be appreciated that additional components, not shown in FIG. **16**, may also be a part of the data processing system **1600** in certain embodiments of the invention, and in certain embodiments of the invention fewer components than shown in FIG. **16** may be used. In addition, it will be appreciated that one or more buses, not shown in FIG. **16**, may be used to interconnect the various components as is well known in the art.

The memory **1610** may store data and/or programs for execution by the data processing system **1600**. The audio input/output **1640** may include a microphone and/or a speaker to, for example, play music and/or provide telephony functionality through the speaker and microphone. The display controller and display device **1660** may include a graphical user interface (GUI). The wireless (e.g., RF) transceivers **1630** (e.g., a WiFi transceiver, an infrared transceiver, a Bluetooth transceiver, a wireless cellular telephony transceiver, etc.) may be used to communicate with other data processing systems. The one or more input devices **1670** allow a user to provide input to the system. These input devices may be a keypad, keyboard, touch panel, multi touch panel, etc. The optional other input/output **1650** may be a connector for a dock.

Embodiments of the invention may include various steps as set forth above. The steps may be embodied in machine-executable instructions which cause a general-purpose or special-purpose processor to perform certain steps. Alternatively, these steps may be performed by specific hardware components that contain hardwired logic for performing the steps, or by any combination of programmed computer components and custom hardware components.

Elements of the present invention may also be provided as a machine-readable medium for storing the machine-executable program code. The machine-readable medium may include, but is not limited to, floppy diskettes, optical disks, CD-ROMs, and magneto-optical disks, ROMs, RAMs, EPROMs, EEPROMs, magnetic or optical cards, or other type of media/machine-readable medium suitable for storing electronic program code.

Throughout the foregoing description, for the purposes of explanation, numerous specific details were set forth in order to provide a thorough understanding of the invention. It will be apparent, however, to one skilled in the art that the invention may be practiced without some of these specific details. For example, it will be readily apparent to those of skill in the art that the functional modules and methods described herein may be implemented as software, hardware or any combination thereof. Moreover, although some embodiments of the invention are described herein within the context of a mobile computing environment, the underlying principles of the invention are not limited to a mobile computing implementation. Virtually any type of client or peer data processing

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devices may be used in some embodiments including, for example, desktop or workstation computers. Accordingly, the scope and spirit of the invention should be judged in terms of the claims which follow.

We claim:

1. A machine-implemented method for authenticating a user over a network comprising:

transmitting a query for client information from an authentication server to a client device, the client information including information related to authentication devices coupled to the client device;

analyzing, by a policy filter at the client device, the query to determine an appropriate privacy class to be used for providing the client information to the authentication server, wherein a privacy class is defined based on a risk that the client information uniquely identifies a user or client device;

wherein the policy filter determines a subset of the client information based on the determined privacy class, where each privacy class specifies a different subset of the client information usable to identify the user and authentication device information usable to identify the client device, where the privacy class is selected from a group consisting of multiple privacy levels such that higher privacy levels divulge less information about the user or the client device to perform authentication;

providing, by the client device, the subset of the client information to the authentication server, the subset of the client information including the information related to the authentication devices coupled to the client device; and

using the subset of the client information within an authentication framework to provide user authentication services over the network.

2. The method as in claim 1 wherein the privacy class is selected from a group consisting of a first privacy class providing a first level of privacy, a second privacy class providing a second level of privacy, and a third privacy class providing a third level of privacy.

3. The method as in claim 2 wherein the third privacy class is associated with a 100% likelihood of the user or machine being able to be uniquely identified, the second privacy class is associated with a second likelihood less than 100%, and the first privacy class is associated with a first likelihood less than the second likelihood.

4. The method as in claim 1 wherein the query for client information sent from the authentication server includes an attribute specifying the privacy class of information that needs to be retrieved.

5. The method as in claim 1 wherein a separate privacy class is selected for each relying party.

6. The method as in claim 5 further comprising: associating information for a highest approved privacy class for each relying party within a secure storage.

7. The method as in claim 6 wherein, whenever the relying party asks for information belonging to a higher privacy class than the one already associated with the relying party, the user is prompted to permanently approve or reject this new privacy class for this relying party.

8. The method as in claim 1 wherein the subset of client information selected based on the determined privacy class includes one or more of a machine model identifier (MMID), client software information, client capabilities, and the information related to each authentication device coupled to the client-device.

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9. The method as in claim 8 wherein the information related to each authentication device coupled to the client device includes one or more of a device ID code, vendor ID code, and device class ID.

10. The method as in claim 1 wherein the authentication server uses the subset of client information for one or more subsequent transactions with the client device.

11. The method as in claim 10 wherein the one or more subsequent transactions include authenticating the user over the network.

12. A system comprising:

an authentication server to transmit a query for client information to a client device, the client information including information related to authentication devices coupled to the client device;

analyzing, by a policy filter at the client device, the query to determine an appropriate privacy class to be used for providing the client information to the authentication server, wherein a privacy class is defined based on a risk that the client information uniquely identifies a user or client device;

wherein the policy filter determines a subset of the client information based on the determined privacy class, where each privacy class specifies a different subset of the client information usable to identify the user and authentication device information usable to identify the client device, where the privacy class is selected from a group consisting of multiple privacy levels such that higher privacy levels divulge less information about the user or the client device to perform authentication;

the client device providing the subset of the client information to the authentication server, the subset of the client information including the information related to the authentication devices coupled to the client device; and

the authentication server using the subset of the client information within an authentication framework to provide user authentication services over a network.

13. The system as in claim 12 wherein the privacy class is selected from a group consisting of a first privacy class providing a first level of privacy, a second privacy class providing a second level of privacy, and a third privacy class providing a third level of privacy.

14. The system as in claim 13 wherein the third privacy class is associated with a 100% likelihood of the user or machine being able to be uniquely identified, the second privacy class is associated with a second likelihood less than 100%, and the first privacy class is associated with a first likelihood less than the second likelihood.

15. The system as in claim 12 wherein the query for client information sent from the authentication server includes an attribute specifying the privacy class of information that needs to be retrieved.

16. The system as in claim 12 wherein a separate privacy class is selected for each relying party.

17. The system as in claim 16 further comprising: the client associating information for a highest approved privacy class for each relying party within a secure storage.

18. The system as in claim 17 wherein, whenever the relying party asks for information belonging to a higher privacy class than the one already associated with the relying party, the user is prompted to permanently approve or reject this new privacy class for this relying party.

19. The system as in claim 12 wherein the subset of client information selected based on the determined privacy class includes one or more of a machine model identifier (MMID),



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client software information, client capabilities, and the information related to each authentication device coupled to the client-device.

20. The system as in claim 19 wherein the information related to each authentication device coupled to the client device includes one or more of a device ID code, vendor ID code, and device class ID.

21. The system as in claim 12 wherein the authentication server uses the subset of client information for one or more subsequent transactions with the client.

22. The system as in claim 21 wherein the one or more subsequent transactions include authenticating a user over a network.

23. A non-transitory machine-readable medium having program code stored thereon which, when executed by a machine, causes the machine to perform the operations of:

transmitting a query for client information from an authentication server to a client, the client information including information related to authentication devices coupled to the client;

analyzing, by a policy filter at the client, the query to determine an appropriate privacy class to be used for providing the client information to the authentication server, wherein a privacy class is defined based on a risk that the client information uniquely identifies a user or client device;

wherein the policy filter determines a subset of the client information based on the determined privacy class, where each privacy class specifies a different subset of the client information usable to identify the user and authentication device information usable to identify the client device, where the privacy class is selected from a group consisting of multiple privacy levels such that higher privacy levels divulge less information about the user or the client device to perform authentication;

providing, by the client, the subset of the client information to the authentication server class, the subset of the client information including the information related to the authentication devices coupled to the client; and

using the subset of the client information within an authentication framework to provide user authentication services over a network.

24. The non-transitory machine-readable medium as in claim 23 wherein the privacy class is selected from a group consisting of a first privacy class providing a first level of

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privacy, a second privacy class providing a second level of privacy, and a third privacy class providing a third level of privacy.

25. The non-transitory machine-readable medium as in claim 24 wherein the third privacy class is associated with a 100% likelihood of the user or machine being able to be uniquely identified, the second privacy class is associated with a second likelihood less than 100%, and the first privacy class is associated with a first likelihood less than the second likelihood.

26. The non-transitory machine-readable medium as in claim 23 wherein the query for client information sent from the authentication server includes an attribute specifying the privacy class of information that needs to be retrieved.

27. The non-transitory machine-readable medium as in claim 23 wherein a separate privacy class is selected for each relying party.

28. The non-transitory machine-readable medium as in claim 27 comprising additional program code to cause the machine to perform the operations of:

associating information for a highest approved privacy class for each relying party within a secure storage.

29. The non-transitory machine-readable medium as in claim 28 wherein, whenever the relying party asks for information belonging to a higher privacy class than the one already associated with the relying party, the user is prompted to permanently approve or reject this new privacy class for this relying party.

30. The non-transitory machine-readable medium as in claim 23 wherein the subset of client information selected based on the determined privacy class includes one or more of a machine model identifier (MMID), client software information, client capabilities, and the information related to each authentication device coupled to the client device.

31. The non-transitory machine-readable medium as in claim 30 wherein the information related to each authentication device coupled to the client includes one or more of a device ID code, vendor ID code, and device class ID.

32. The non-transitory machine-readable medium as in claim 23 wherein the authentication server uses the subset of client information for one or more subsequent transactions with the client.

33. The non-transitory machine-readable medium as in claim 32 wherein the one or more subsequent transactions include authenticating a user over a network.

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